

Address-of and dereference operators

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& address-of operator

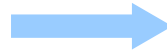
***** dereference operator

Address-of operator and dereferencing operator

address-of operator **&** : the address of a variable

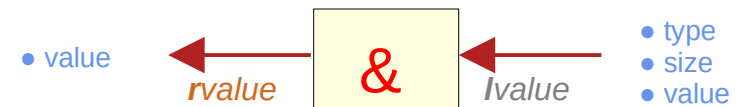
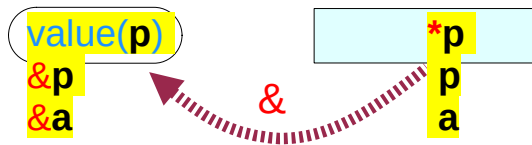
a **variable** has a memory location
whose value can be changed
by an assignment

a **variable** must be an **lvalue**



&variable returns
the address of a variable

&variable returns an **rvalue**



Address-of operator and dereferencing operator

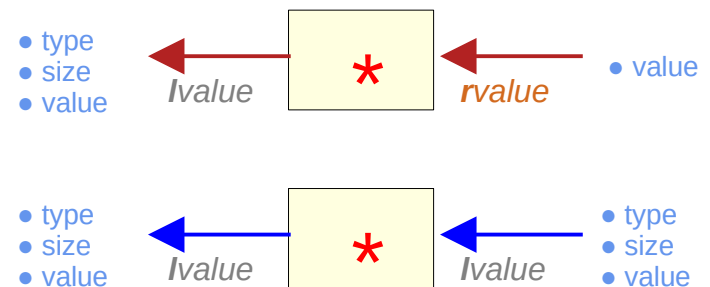
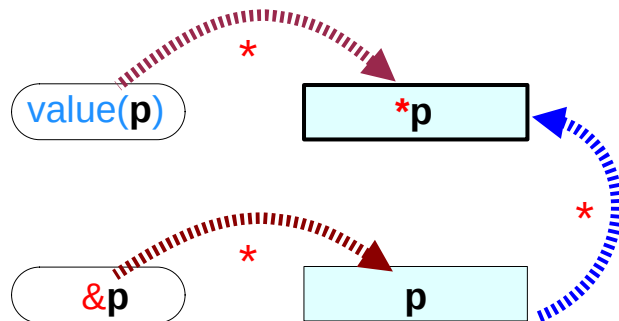
*dereferencing operator ***** : the content at an address*

*an **address** can be an address value or a pointer variable that holds an address value*

*an **address** must be an **rvalue**, or evaluated to be an **rvalue***

**** address** returns the content value at the address*

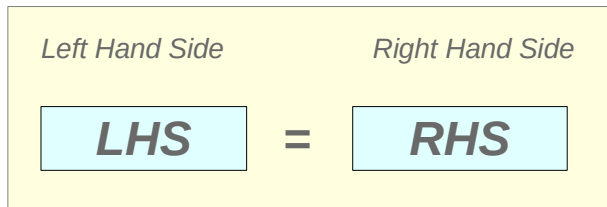
**** address** returns an **lvalue***



***p** $\equiv$ ***value(p)**

Ivalue and rvalue in assignments

an assignment statement



- in the **LHS**, only **Ivalue** can exist
- **rvalue** can exist only in the **RHS**

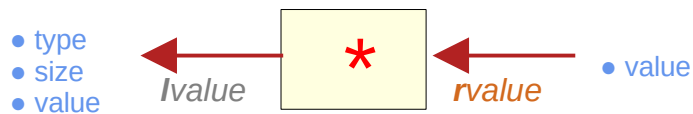
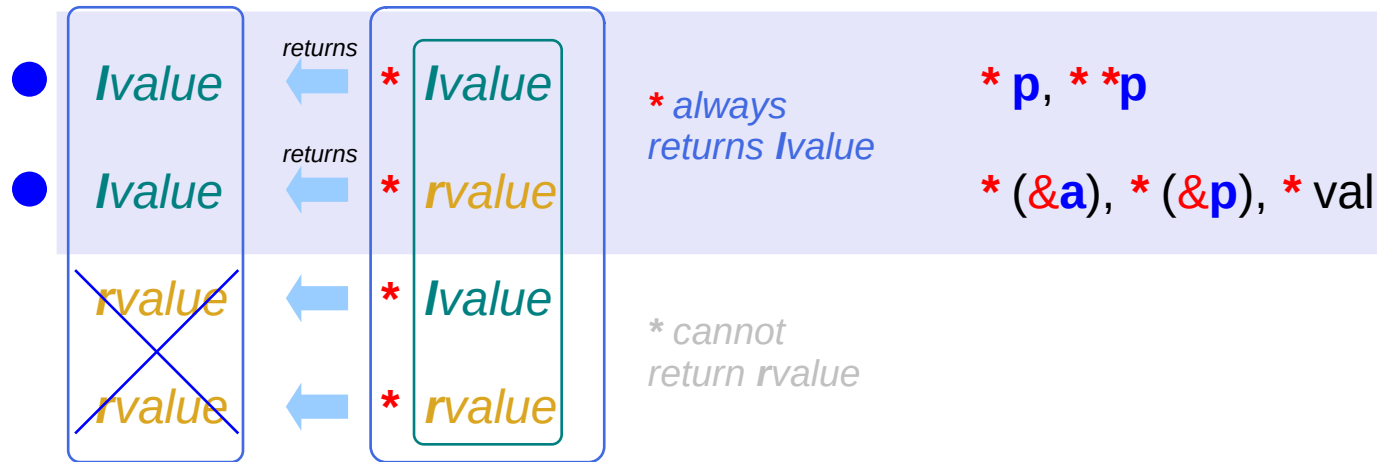
```
int  a, b = 10 ;  
int  *p, *q = &a ;
```

a, b, p, q	: Ivalues	... variables	... RW
*p, *q	: Ivalues	... variables	... RW
&a, &b	: rvalues	... constants	... RO

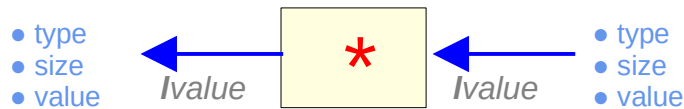
Ivalue	=	Ivalue
Ivalue	=	rvalue
rvalue	=	Ivalue
rvalue	=	rvalue

p	=	q ;
p	=	&a ;
&a	=	p ;
&a	=	&b ;

Ivalue and rvalue with * and & operators



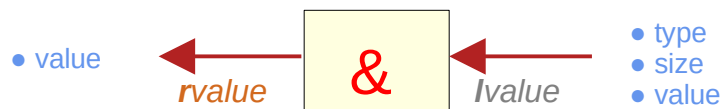
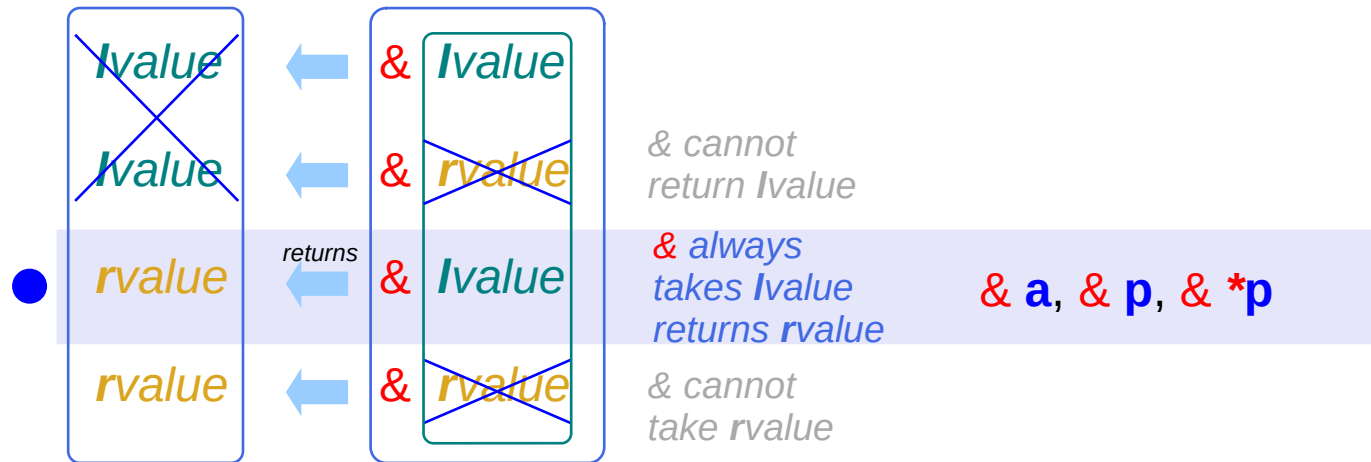
* can be applied to either an **Ivalue** variable or a **rvalue** address



* **operand** becomes an **Ivalue** variable thus it can be applied successively.

$$*p \equiv *value(p)$$

Ivalue and rvalue with * and & operators



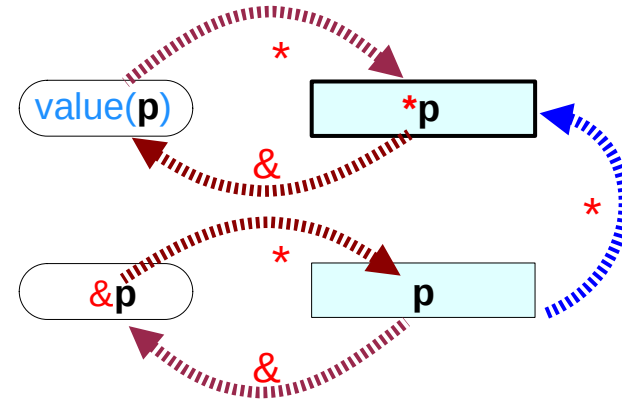
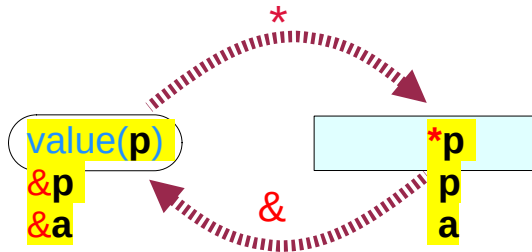
& can be applied
to only an **lvalue** variable and
returns only an **rvalue** address

Variable types

a	non-pointer variables	<i>lvalue</i>
p	pointer variables	<i>lvalue</i>
*p	dereferenced variables	<i>lvalue</i>
val	address values	<i>rvalue</i>

C operators : *, &

C operators : *, &

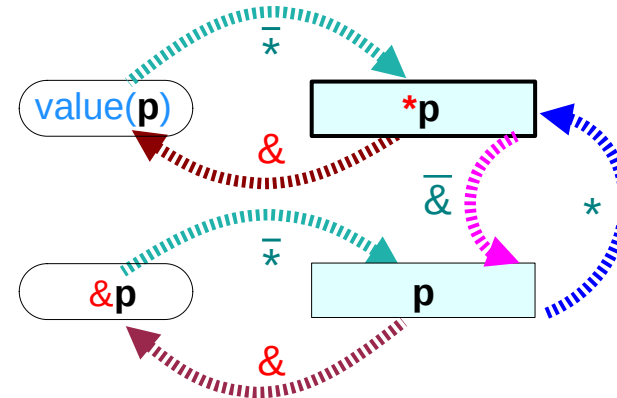
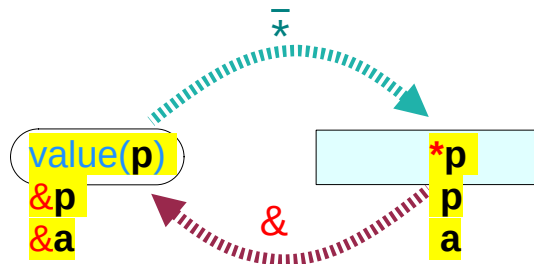


	*	a	<i>Ivalue</i>
<i>Ivalue</i> p ←	*	p	<i>Ivalue</i>
<i>Ivalue</i> * p ←	*	* p	<i>Ivalue</i>
<i>Ivalue</i> * val ←	*	val	<i>rvalue</i>

<i>rvalue</i> & p ←	&	a	<i>Ivalue</i>
<i>rvalue</i> & p ←	&	p	<i>Ivalue</i>
<i>rvalue</i> & * p ←	&	* p	<i>Ivalue</i>
	&	val	<i>rvalue</i>

Extending C operators with $\bar{*}$, $\bar{\&}$ (1)

Extending C operators with Mathematical operators ($\bar{*}$, $\bar{\&}$)



		*	a	Ivalue
Ivalue	p ←	*	p	Ivalue
Ivalue	*p ←	*	*p	Ivalue
★ Ivalue	$\bar{*}$ val ←	$\bar{*}$	val	rvalue

rvalue	& p ←	&	a	Ivalue
rvalue	& p ←	&	p	Ivalue
rvalue	& *p ←	&	*p	Ivalue
★ Ivalue	p ←	$\bar{\&}$	*p	Ivalue

Extending C operators with $\bar{*}$, $\bar{\&}$ (2)

Extending C operators with Mathematical operators ($\bar{*}$, $\bar{\&}$)

	*	a	Ivalue
Ivalue p ←	*	p	Ivalue
Ivalue * p ←	*	* p	Ivalue
★	*	val	rvalue

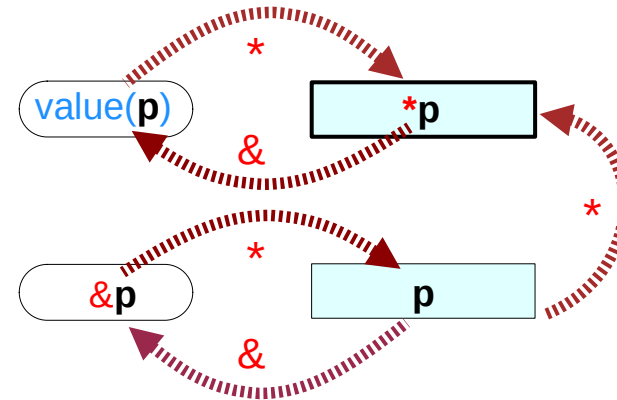
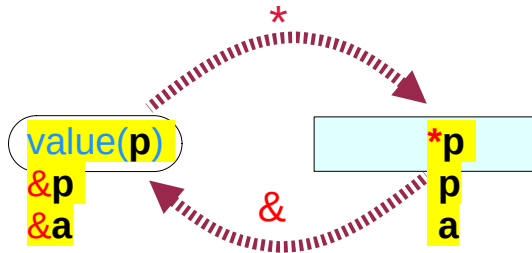
rvalue & p ←	&	a	Ivalue
rvalue & p ←	&	p	Ivalue
rvalue & * p ←	&	* p	Ivalue
	&	val	rvalue

	$\bar{*}$	a	Ivalue
	$\bar{*}$	p	Ivalue
	$\bar{*}$	* p	Ivalue
★ Ivalue $\bar{*}$ val ←	$\bar{*}$	val	rvalue

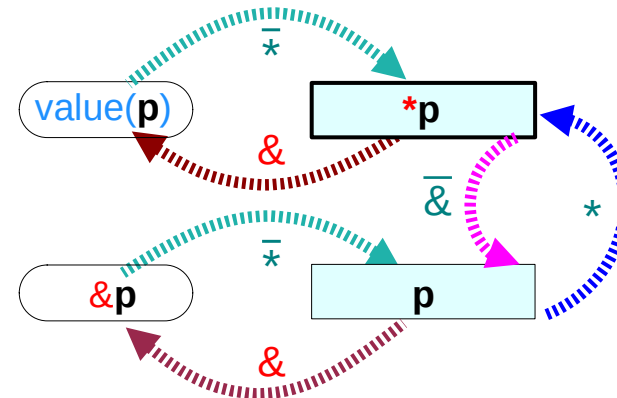
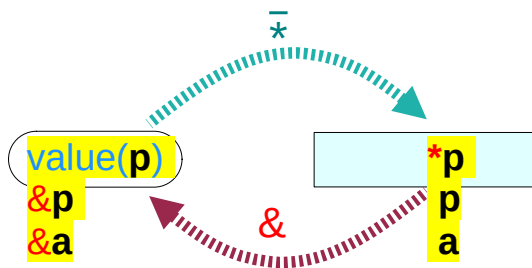
	$\bar{\&}$	a	Ivalue
	$\bar{\&}$	p	Ivalue
★ Ivalue p ←	$\bar{\&}$	* p	Ivalue
	$\bar{\&}$	val	rvalue

Comparisons (1)

C operators : $*$, $\&$



Extending C operators with Mathematical operators ($\bar{*}$, $\bar{\&}$)



Comparisons (2)

C operators : *, &

	*	a	Ivalue
Ivalue p ←	*	p	Ivalue
Ivalue * p ←	*	* p	Ivalue
Ivalue * val ←	*	val	rvalue

rvalue & a ←	&	a	Ivalue
rvalue & p ←	&	p	Ivalue
rvalue & * p ←	&	* p	Ivalue
	&	val	rvalue

Extending C operators with Mathematical operators ($\bar{*}$, $\bar{\&}$)

	*	a	Ivalue
Ivalue p ←	*	p	Ivalue
Ivalue * p ←	*	* p	Ivalue
★ Ivalue $\bar{*}$ val ←	$\bar{*}$	val	rvalue

rvalue & a ←	&	a	Ivalue
rvalue & p ←	&	p	Ivalue
rvalue & * p ←	&	* p	Ivalue
★ Ivalue p ←	$\bar{\&}$	* p	Ivalue

Comparisons (3)

C operators : *, &

	*	a	<i>Ivalue</i>
<i>Ivalue</i> p ←	*	p	<i>Ivalue</i>
<i>Ivalue</i> * p ←	*	* p	<i>Ivalue</i>
<i>Ivalue</i> * val ←	*	val	<i>rvalue</i>

<i>rvalue</i> & p ←	&	a	<i>Ivalue</i>
<i>rvalue</i> & p ←	&	p	<i>Ivalue</i>
<i>rvalue</i> & * p ←	&	* p	<i>Ivalue</i>
	&	val	<i>rvalue</i>

Extending C operators with Mathematical operators ($\bar{*}$, $\bar{\&}$)

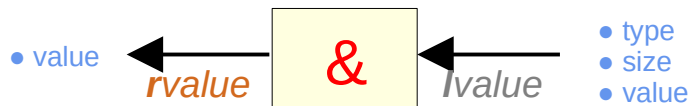
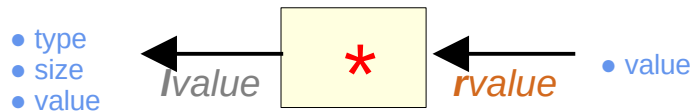
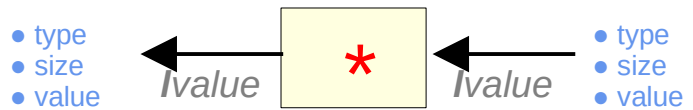
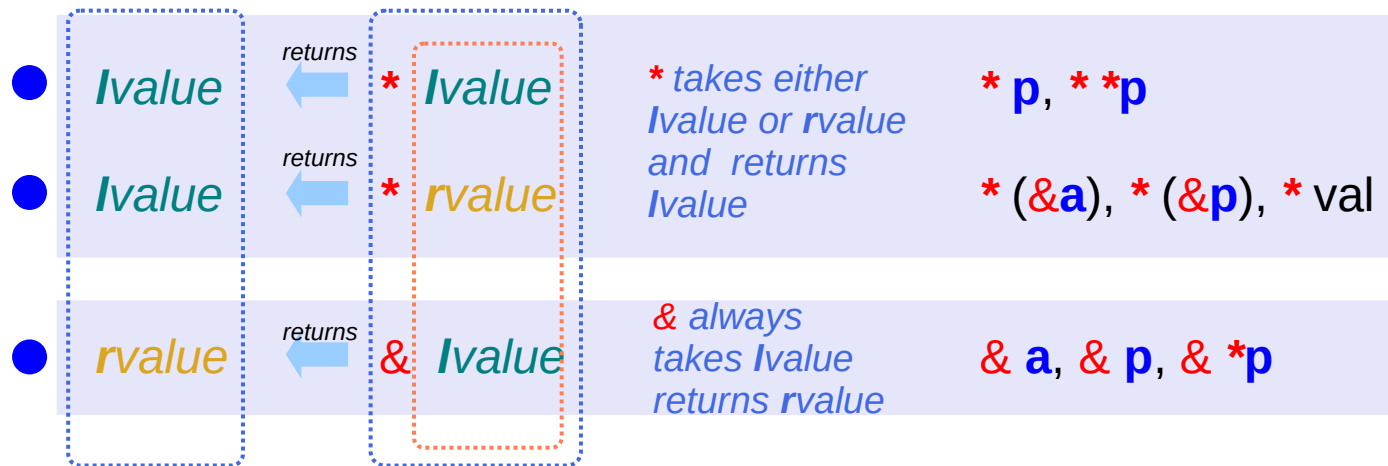
	*	a	<i>Ivalue</i>
<i>Ivalue</i> p ←	*	p	<i>Ivalue</i>
<i>Ivalue</i> * p ←	*	* p	<i>Ivalue</i>
★	*	val	<i>rvalue</i>

<i>rvalue</i> & p ←	&	a	<i>Ivalue</i>
<i>rvalue</i> & p ←	&	p	<i>Ivalue</i>
<i>rvalue</i> & * p ←	&	* p	<i>Ivalue</i>
	&	val	<i>rvalue</i>

	$\bar{*}$	a	<i>Ivalue</i>
	$\bar{*}$	p	<i>Ivalue</i>
	$\bar{*}$	* p	<i>Ivalue</i>
★ <i>Ivalue</i> $\bar{*}$ val ←	$\bar{*}$	val	<i>rvalue</i>

	$\bar{\&}$	a	<i>Ivalue</i>
	$\bar{\&}$	p	<i>Ivalue</i>
★ <i>Ivalue</i> p ←	$\bar{\&}$	* p	<i>Ivalue</i>
	$\bar{\&}$	val	<i>rvalue</i>

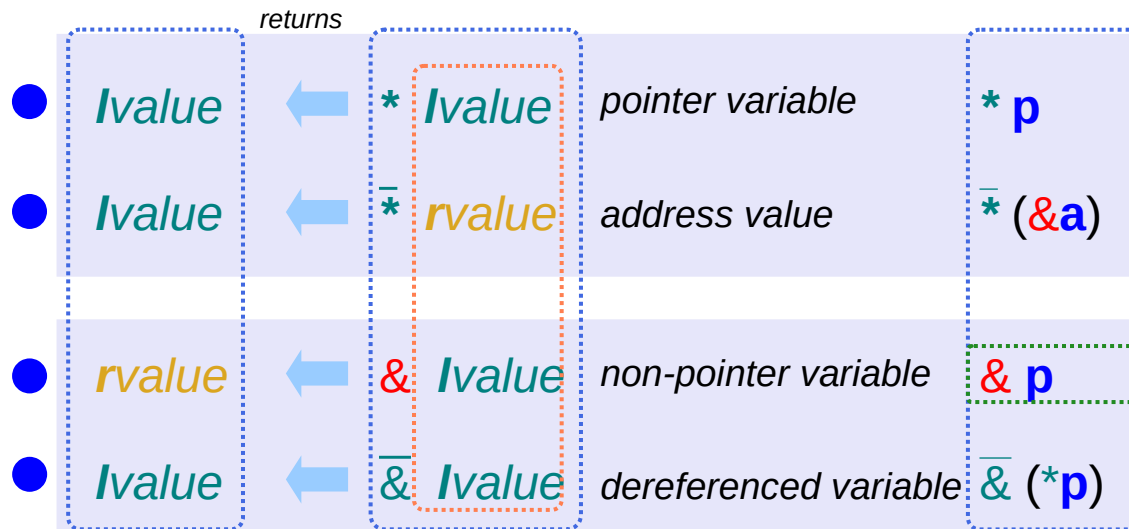
Ivalue and rvalue with * and & operators



- * (an Ivalue variable)
- * (an rvalue address)
- * **operand** : an Ivalue variable
- * can be applied successively.

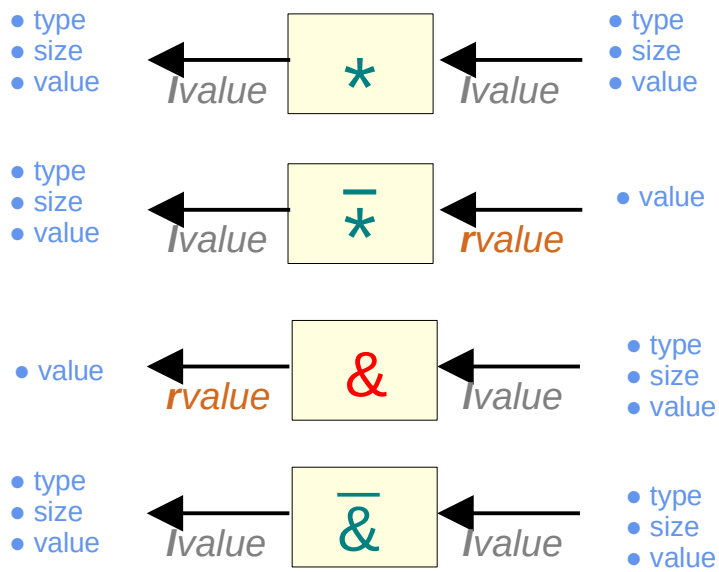
- & (an Ivalue variable)
- & **operand** : an rvalue address,
- & cannot be applied successively.

Ivalue and rvalue with * and & operators



* can be applied successively.
 $\overline{\&}$ can be applied successively.

****p**
 $\overline{\&} **p = *p$
 $\overline{\&} \& **p = \overline{\&} *p = p$



Recursive application of the address-of operator

$$* = \overline{*} \overline{v}$$

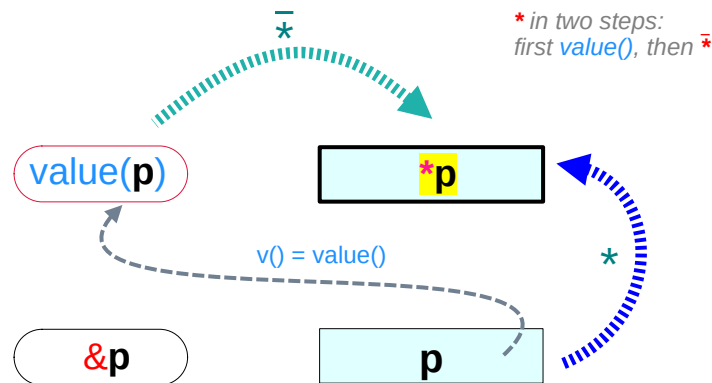
$$*v^{-1} = \overline{*} \overline{vv^{-1}}$$

$$*v^{-1} = \overline{*}$$

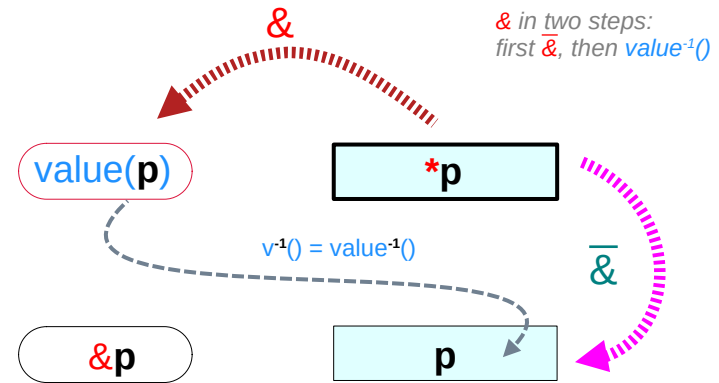
$$\overline{\&} = v^{-1} \&$$

$$v \overline{\&} = vv^{-1} \&$$

$$v \overline{\&} = \&$$



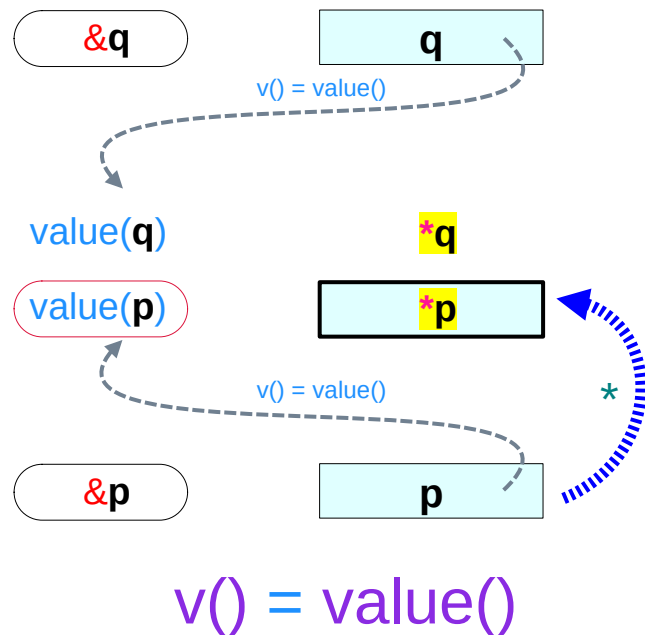
$$v() = \text{value}()$$



$$v^{-1}() = \text{value}^{-1}()$$

When two pointers point to the same location

When two pointers **p** and **q**
point to the same memory location
 $\text{value}(\mathbf{p}) = \text{value}(\mathbf{q})$
 $\text{value}^{-1}(*\mathbf{p}) = \mathbf{p}$ or
 $\text{value}^{-1}(*\mathbf{q}) = \mathbf{q}$
Cannot determine



$\text{value}(\mathbf{p}) = \text{value}(\mathbf{q}) = \mathbf{x}$
then $\text{value}^{-1}(\mathbf{x}) = \mathbf{p}$ or $\mathbf{q}$?

When $\text{value}^{-1}()$ is prohibited (1)

$$* = \bar{*}v$$

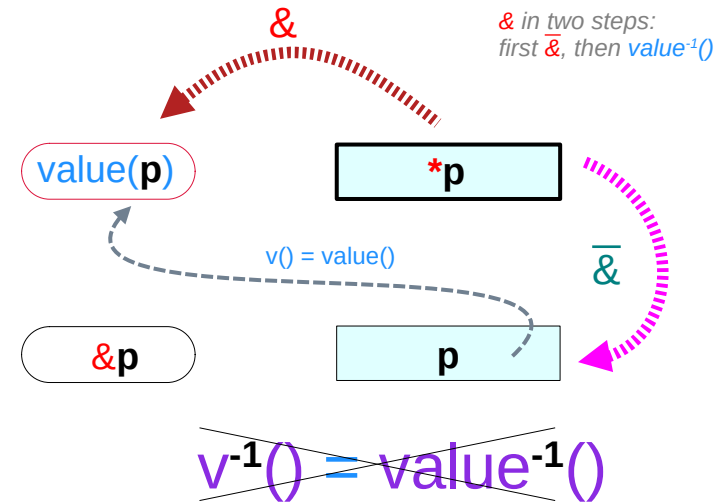
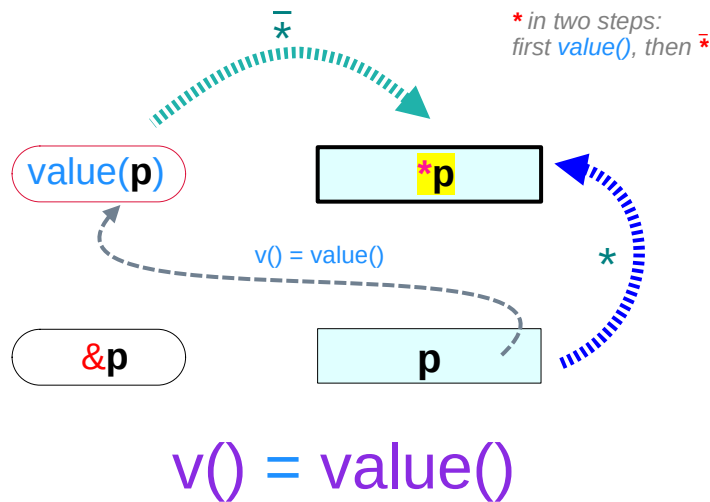
~~$$*v^{-1} = \bar{*}vv^{-1}$$~~

~~$$*v^{-1} = \bar{*}$$~~

~~$$\bar{\&} = v^{-1}\&$$~~

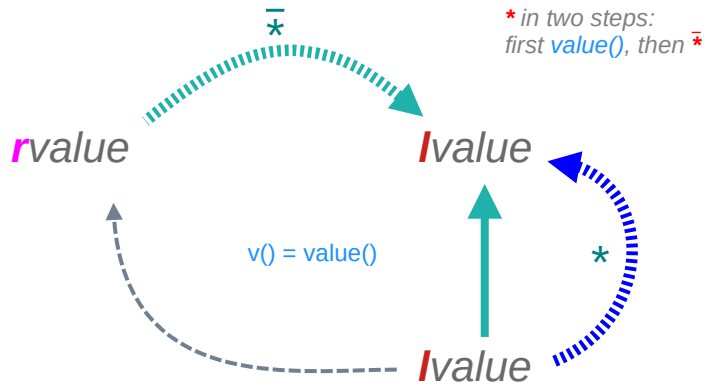
~~$$v\bar{\&} = vv^{-1}\&$$~~

$$v\bar{\&} = \&$$



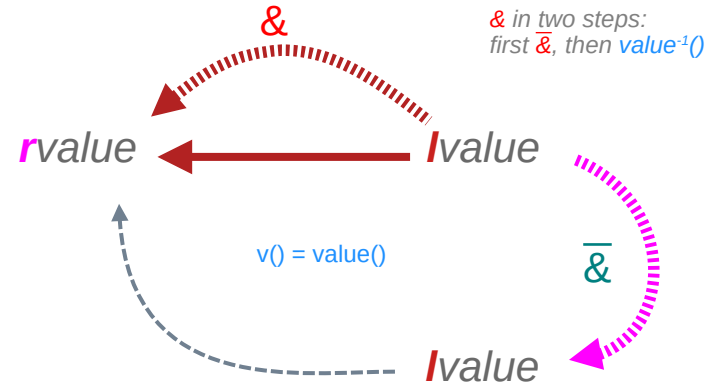
When $\text{value}^{-1}()$ is prohibited (2)

$$* = \overline{*} v$$



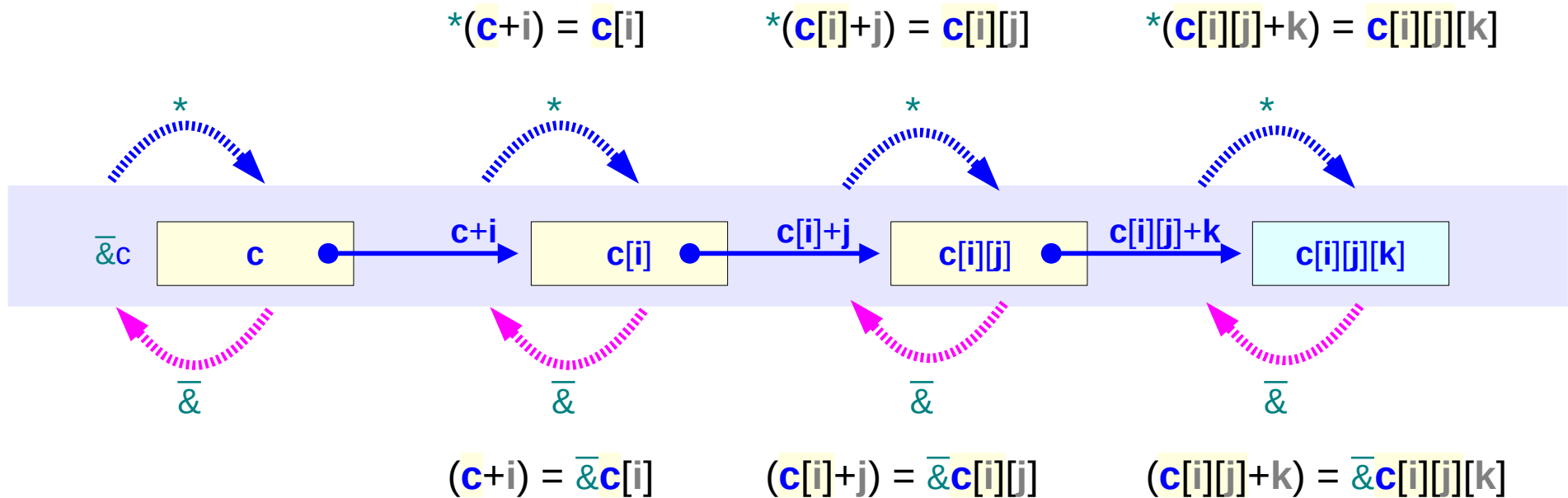
$$v() = \text{value}()$$

$$v \overline{\&} = \&$$

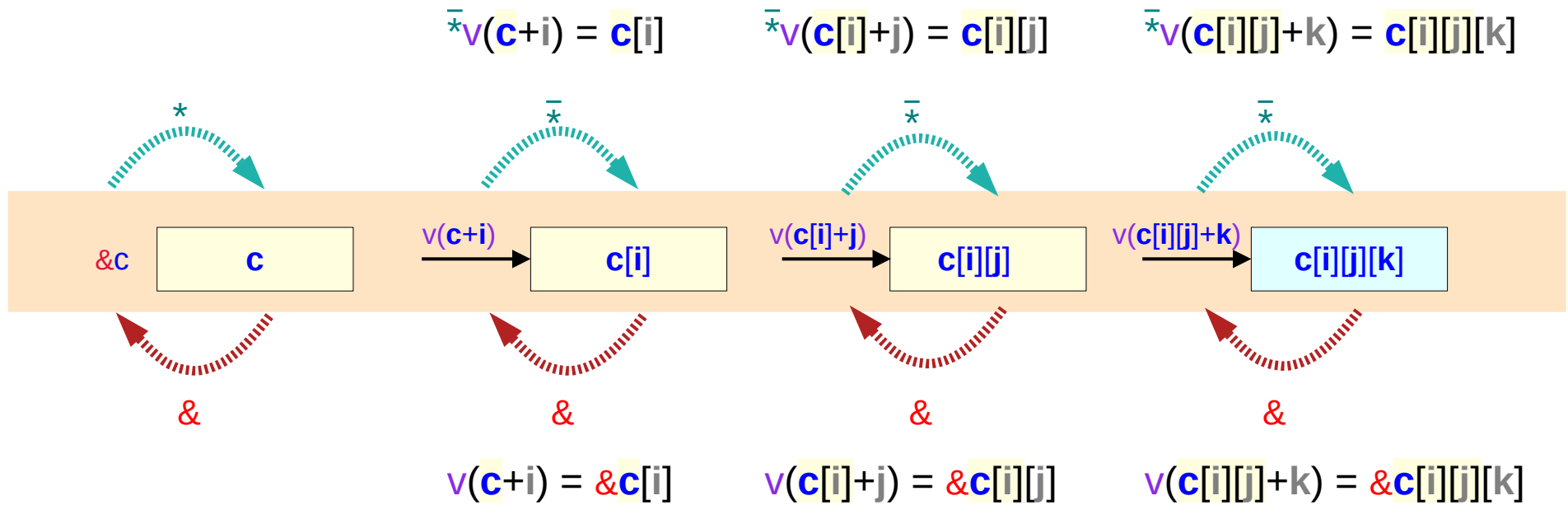


$$v() = \text{value}()$$

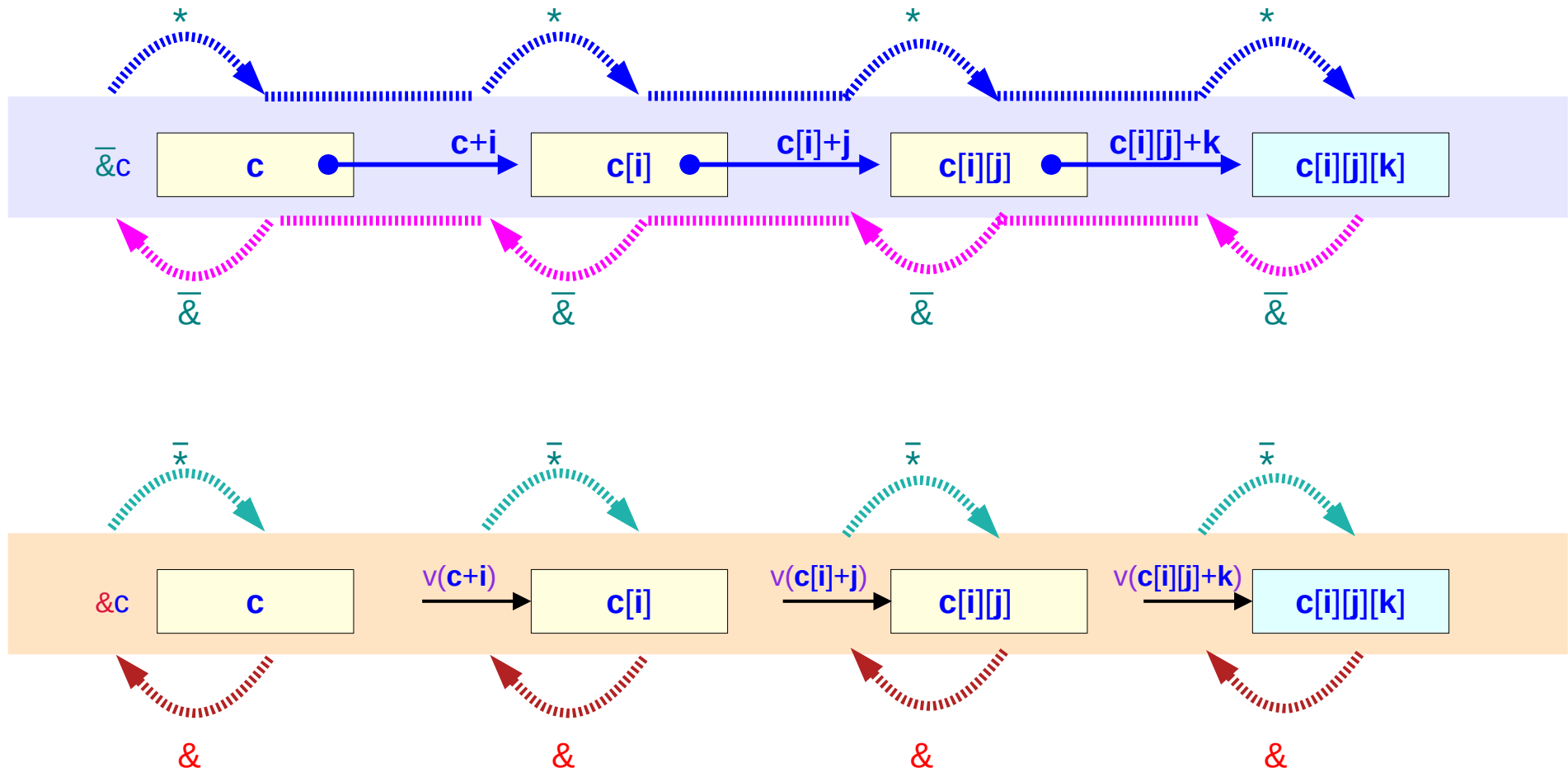
Inverse relationship between $*$ and $\bar{\&}$ operators



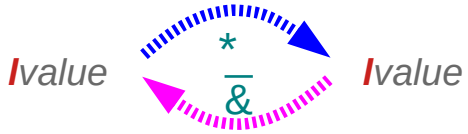
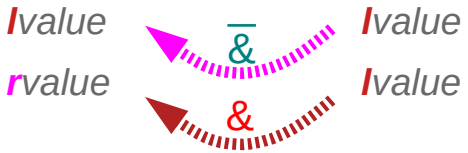
Inverse relationship between $\bar{*}$ and $\&$ operators



Inverse relationship (1)



Inverse relationship (2)

 Diagram: A blue arrow labeled '*' points from an 'lvalue' to another 'lvalue'. A pink arrow labeled '&' points from an 'lvalue' to another 'lvalue'.	$*(c+i) = c[i]$ $(c+i) = \bar{\&c}[i]$	$*(c[i]+j) = c[i][j]$ $(c[i]+j) = \bar{\&c}[i][j]$	$*(c[i][j]+k) = c[i][j][k]$ $(c[i][j]+k) = \bar{\&c}[i][j][k]$
 Diagram: A green arrow labeled '*' points from an 'rvalue' to an 'lvalue'. A red arrow labeled '&' points from an 'lvalue' to an 'rvalue'.	$\bar{*}v(c+i) = c[i]$ $v(c+i) = \&c[i]$	$\bar{*}v(c[i]+j) = c[i][j]$ $v(c[i]+j) = \&c[i][j]$	$\bar{*}v(c[i][j]+k) = c[i][j][k]$ $v(c[i][j]+k) = \&c[i][j][k]$
 Diagram: A blue arrow labeled '*' points from an 'lvalue' to an 'rvalue'. A green arrow labeled '*' points from an 'rvalue' to an 'lvalue'.	$*(c+i) = c[i]$ $\bar{*}v(c+i) = c[i]$	$*(c[i]+j) = c[i][j]$ $\bar{*}v(c[i]+j) = c[i][j]$	$*(c[i][j]+k) = c[i][j][k]$ $\bar{*}v(c[i][j]+k) = c[i][j][k]$
 Diagram: A pink arrow labeled '&' points from an 'lvalue' to an 'lvalue'. A red arrow labeled '&' points from an 'rvalue' to an 'rvalue'.	$(c+i) = \bar{\&c}[i]$ $v(c+i) = \&c[i]$	$(c[i]+j) = \bar{\&c}[i][j]$ $v(c[i]+j) = \&c[i][j]$	$(c[i][j]+k) = \bar{\&c}[i][j][k]$ $v(c[i][j]+k) = \&c[i][j][k]$

Types of sub-arrays (1)

```
s2.c: In function 'main':
s2.c:7:12: warning: format '%d' expects argument of type 'int', but argument 2 has type 'int (*)[3][4]' [-Wformat=]
 7 | printf("%d\n", a);
   |      ^~
   |      | |
   |      int int (*)[3][4]
s2.c:8:12: warning: format '%d' expects argument of type 'int', but argument 2 has type 'int (*)[4]' [-Wformat=]
 8 | printf("%d\n", a[0]);
   |      ^~
   |      | |
   |      int int (*)[4]
s2.c:9:12: warning: format '%d' expects argument of type 'int', but argument 2 has type 'int *' [-Wformat=]
 9 | printf("%d\n", a[0][0]);
   |      ^~
   |      | |
   |      int int *
   |      %ls

s2.c:11:12: warning: format '%d' expects argument of type 'int', but argument 2 has type 'int (*)[3][4]' [-Wformat=]
11 | printf("%d\n", a+1);
   |      ^~
   |      | |
   |      int int (*)[3][4]
s2.c:12:12: warning: format '%d' expects argument of type 'int', but argument 2 has type 'int (*)[4]' [-Wformat=]
12 | printf("%d\n", a[0]+1);
   |      ^~
   |      | |
   |      int int (*)[4]
s2.c:13:12: warning: format '%d' expects argument of type 'int', but argument 2 has type 'int *' [-Wformat=]
13 | printf("%d\n", a[0][0]+1);
   |      ^~
   |      | |
   |      int int *
```

```
#include <stdio.h>

int main(void) {

    int a[2][3][4];

    printf("%d\n", a);
    printf("%d\n", a[0]);
    printf("%d\n", a[0][0]);

    printf("%d\n", a+1);
    printf("%d\n", a[0]+1);
    printf("%d\n", a[0][0]+1);

    printf("%d\n", &a[0]);
    printf("%d\n", &a[0][0]);
    printf("%d\n", &a[0][0][0]);

}
```

Types of sub-arrays (2)

```
s2.c:15:12: warning: format '%d' expects argument of type 'int', but argument 2 has type 'int (*)[3][4]' [-Wformat=]
15 | printf("%d\n", &a[0]);
   |           ^~  ~~~~~
   |           |  |
   |           int int (*)[3][4]
s2.c:16:12: warning: format '%d' expects argument of type 'int', but argument 2 has type 'int (*)[4]' [-Wformat=]
16 | printf("%d\n", &a[0][0]);
   |           ^~  ~~~~~
   |           |  |
   |           int int (*)[4]
s2.c:17:12: warning: format '%d' expects argument of type 'int', but argument 2 has type 'int *' [-Wformat=]
17 | printf("%d\n", &a[0][0][0]);
   |           ^~  ~~~~~
   |           |  |
   |           int int *
   |           %ls
y
```

```
#include <stdio.h>

int main(void) {

    int a[2][3][4];

    printf("%d\n", a);
    printf("%d\n", a[0]);
    printf("%d\n", a[0][0]);

    printf("%d\n", a+1);
    printf("%d\n", a[0]+1);
    printf("%d\n", a[0][0]+1);

    printf("%d\n", &a[0]);
    printf("%d\n", &a[0][0]);
    printf("%d\n", &a[0][0][0]);

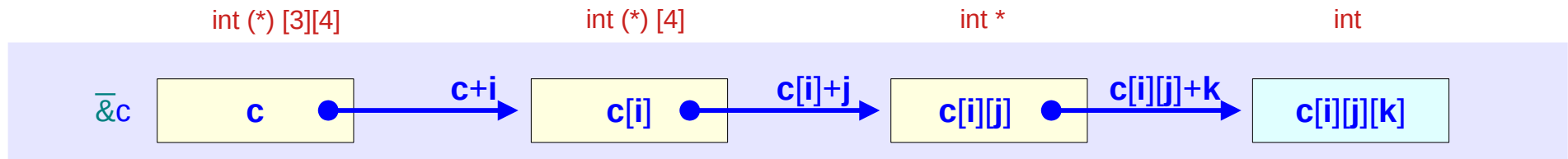
}
```

Types of sub-arrays (3)

```
int c[2][3][4];
```

c		:: int (*) [3][4]
c[i]	= *(c +i)	:: int (*) [4]
c[i][j]	= *(c [i]+j)	:: int *
c[i][j][k]	= *(c [i][j]+k)	:: int

&c[i]	= (c +i)	:: int (*) [3][4]
&c[i][j]	= (c [i]+j)	:: int (*) [4]
&c[i][j][k]	= (c [i][j]+k)	:: int (*)



Dimensional Parentheses (1)

```
int c[2][3][4];
```

$(c+i) = ?$

$((c+i)+j) = ?$

$((c+i)+j)+k = ?$

$\neq c+i$

$\neq c+i+j$

$\neq c+i+j+k$

$(c+i) = (c+i)_{[3][4]}$

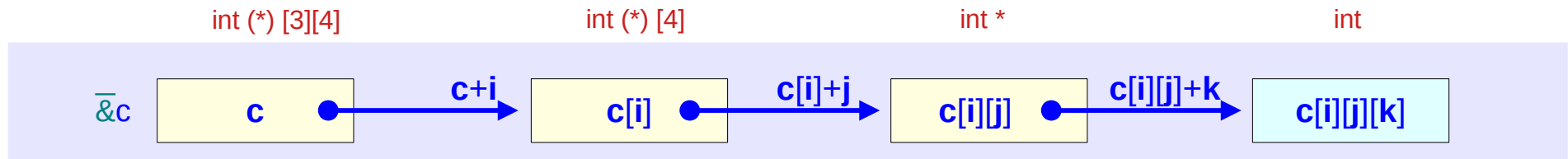
$((c+i)+j) = ((c+i)_{[3][4]}+j)_{[4]}$

$((c+i)+j)+k = (((c+i)_{[3][4]}+j)_{[4]}+k)$

dimensional parentheses

not a simple parenthesis, but associated with the dimensions of an array defined

mathematical expressions obtained by mathematical operator $\bar{\&}$



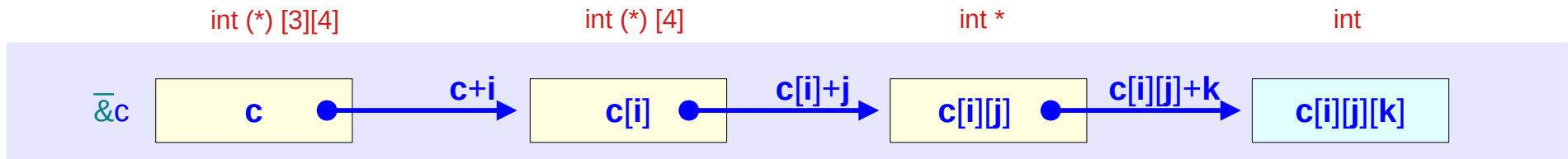
Dimensional Parentheses (2)

```
int c[2][3][4];
```

$(\mathbf{c}+i)$	$= (\mathbf{c}+i)_{[3][4]}$
$((\mathbf{c}+i)+j)$	$= ((\mathbf{c}+i)_{[3][4]}+j)_{[4]}$
$(((\mathbf{c}+i)+j)+k)$	$= (((\mathbf{c}+i)_{[3][4]}+j)_{[4]}+k)$

$\text{type}(\mathbf{c}+i)$	$= \text{int } (*) [3][4]$
$\text{type}((\mathbf{c}+i)+j)$	$= \text{int } (*) [4]$
$\text{type}(((\mathbf{c}+i)+j)+k)$	$= \text{int } *$

$v(\mathbf{c}+i)$	$= v(\mathbf{c}) + i * \text{sizeof}(3*4)$
$v((\mathbf{c}+i)+j)$	$= v(\mathbf{c}+i) + j * \text{sizeof}(4)$
$v(((\mathbf{c}+i)+j)+k)$	$= v((\mathbf{c}+i)+j) + k$



Dimensional Parentheses (3)

```
int c[2][3][4];
```

```
type(c+i)      = int (*) [3][4]
type((c+i)+j)   = int (*) [4]
type(((c+i)+j)+k) = int *
```

```
v(c+i)          = v(c) + i * sizeof(3*4)
v((c+i)+j)       = v(c+i) + j * sizeof(4)
v(((c+i)+j)+k)   = v((c+i)+j) + k
```

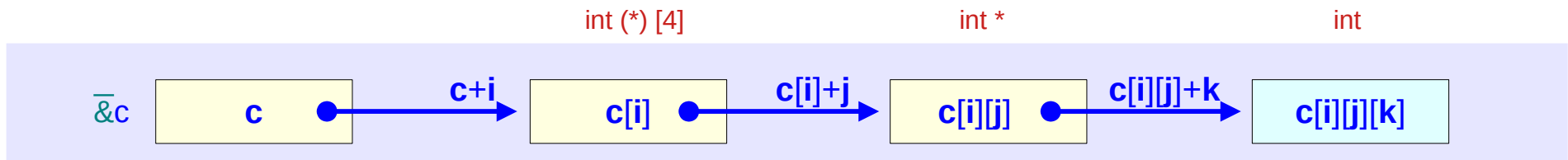
equivalences

```
c[i][j]  ≡ &c[i][j][0]
c[i]     ≡ &c[i][0]
c        ≡ &c[0]
```

address replication

```
value(c[i][j]) = value(&c[i][j])
value(c[i])    = value(&c[i])
value(c)       = value(&c)
```

```
v(c+i)          = v(c[i])
v(c[i]+j)        = v(c[i][j])
v(c[i][j]+k)     = v(c[i][j][k])
```



Address replications in a multi-dimensional array

```
int c [2][3][4] ;
```

equivalences

$c[i][j] \equiv \&c[i][j][0]$
 $c[i] \equiv \&c[i][0]$
 $c \equiv \&c[0]$

address replication

$\text{value}(c[i][j]) = \text{value}(\&c[i][j])$
 $\text{value}(c[i]) = \text{value}(\&c[i])$
 $\text{value}(c) = \text{value}(\&c)$

$c, c[0], c[0][0]$:
these virtual pointers have
the same address value

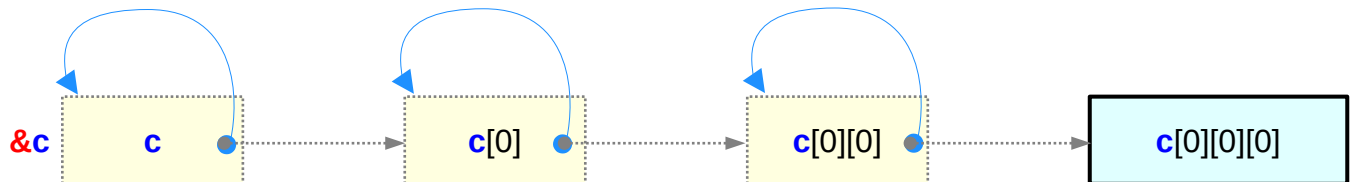
a physical location
has a unique address



$c \equiv \&c[0]$

$c[0] \equiv \&c[0][0]$

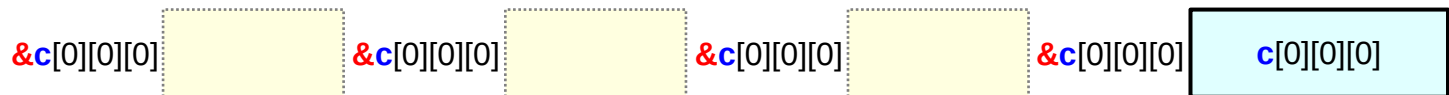
$c[0][0] \equiv \&c[0][0][0]$



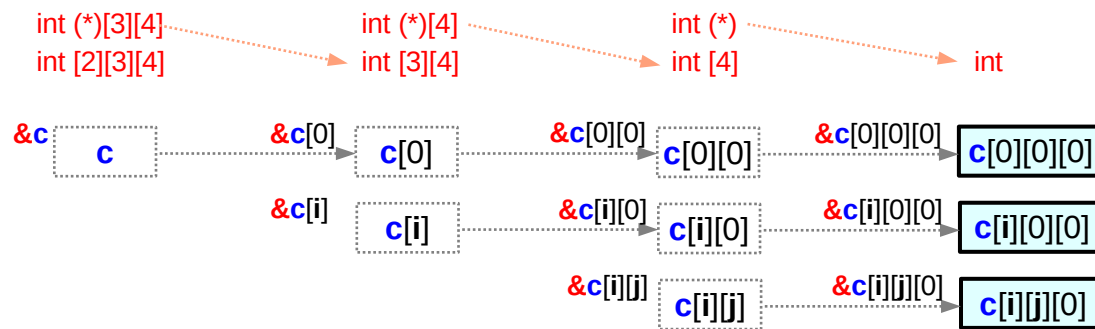
all have the same address value



all have the same starting address



Referencing sub-arrays



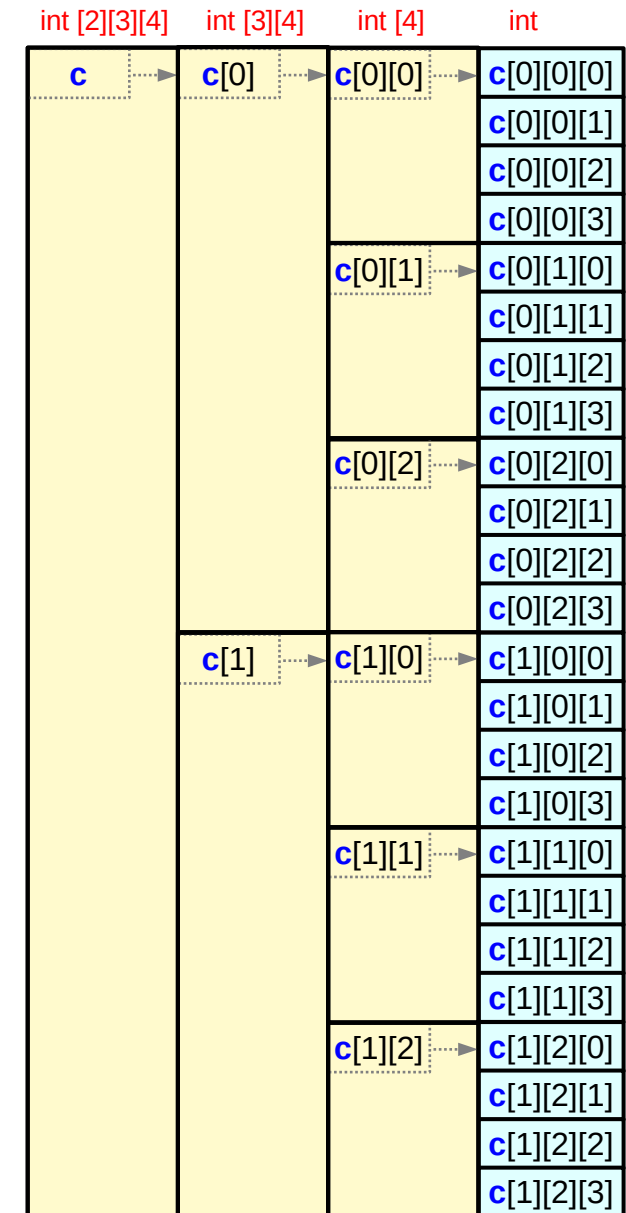
equivalence relations

$c[i][j] \equiv *(c[i]+j)$ $\&c[i][j] \equiv (c[i]+j)$ $\&c[i][0] \equiv c[i]$
 $c[i] \equiv *(c+i)$ $\&c[i] \equiv (c+i)$ $\&c[0] \equiv c$

address replication

$\text{value}(c[i][j]) = \text{value}(\&c[i][j]) = \text{value}(c[i]+j) = * \text{value}(c[i]+j)$
 $\text{value}(c[i]) = \text{value}(\&c[i]) = \text{value}(c+i) = * \text{value}(c+i)$

$c[i], c[i][0]$ point to the same data $c[i][0][0]$
 $c, c[0], c[0][0]$ point to the same data $c[0][0][0]$



Types, sizes, and values of sub-arrays

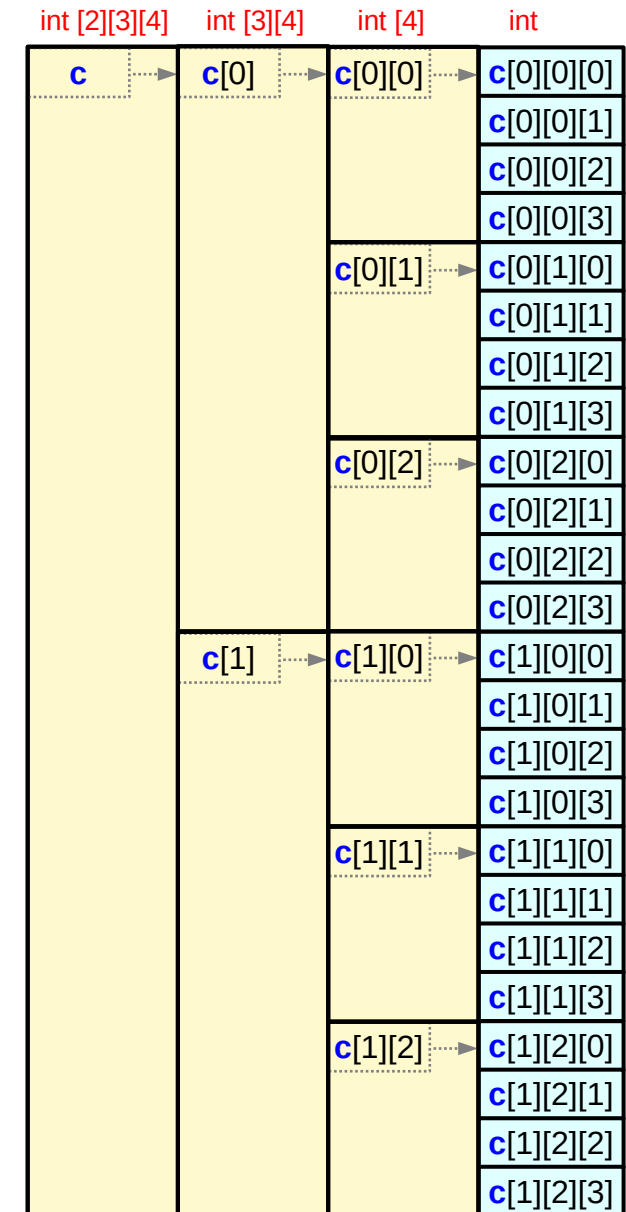
`int c [2][3][4];` static allocation

```
value(c) = value(c[0]) = value(c[0][0]) = &c[0][0][0]
           value(c[0][1]) = &c[0][1][0]
           value(c[0][2]) = &c[0][2][0]
           value(c[1]) = value(c[1][0]) = &c[1][0][0]
           value(c[1][1]) = &c[1][1][0]
           value(c[1][2]) = &c[1][2][0]
```

```
sizeof(c)    = 2*3*4 * sizeof(int)
sizeof(c[i])  = 3*4 * sizeof(int)
sizeof(c[i][j]) = 4 * sizeof(int)
```

```
type(c)      = int [2][3][4]
              int (*)[3][4]
type(c[i])   = int [3][4]
              int (*)[4]
type(c[i][j]) = int [4]
              int (*)
```

pointers to arrays



Dimensional Parentheses

```
int c[2][3][4];
```

$(c+i)$	:: int (*) [3][4]
$((c+i)+j)$	:: int (*) [4]
$((((c+i)+j)+k))$	:: int (*)

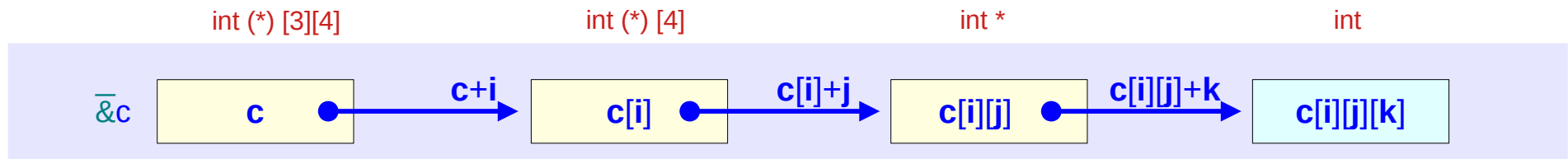
$v(c+i)$	= $v(c) + i * \text{sizeof}(3*4)$
$v((c+i)+j)$	= $v(c+i) + j * \text{sizeof}(4)$
$v(((c+i)+j)+k)$	= $v((c+i)+j) + k$

$(c+i)$	= $(c+i)_{[3][4]}$	$\neq$	$c+i$
$((c+i)+j)$	= $((c+i)_{[3][4]}+j)_{[4]}$	$\neq$	$c+i+j$
$((((c+i)+j)+k))$	= $((((c+i)_{[3][4]}+j)_{[4]}+k))$	$\neq$	$c+i+j+k$

dimensional parentheses

not a simple parenthesis, but associated with the dimensions of an array defined

mathematical expressions obtained by mathematical operator $\bar{\&}$



Types of sub-arrays (2)

```
int c[2][3][4];
```

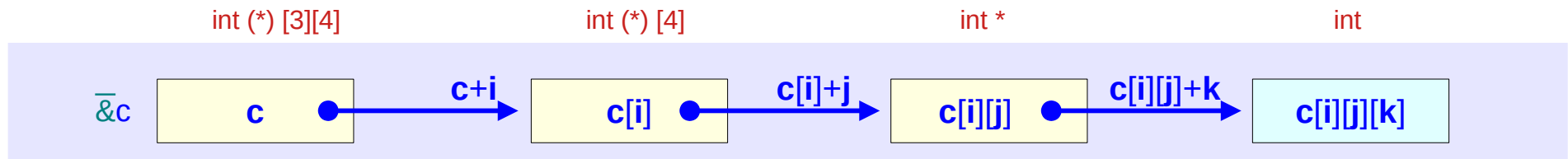
$(c+i)$	:: int (*) [3][4]
$((c+i)+j)$	:: int (*) [4]
$((((c+i)+j)+k))$	:: int (*)

$v(c+i)$	= $v(c) + i * \text{sizeof}(3*4)$
$v((c+i)+j)$	= $v(c+i) + j * \text{sizeof}(4)$
$v(((c+i)+j)+k)$	= $v((c+i)+j) + k$

$(c+i)$	= $(c+i)_{[3][4]}$	$\neq$	$c+i$
$((c+i)+j)$	= $((c+i)_{[3][4]}+j)_{[4]}$	$\neq$	$c+i+j$
$((((c+i)+j)+k))$	= $((((c+i)_{[3][4]}+j)_{[4]}+k))$	$\neq$	$c+i+j+k$

not a simple parenthesis, but associated with the dimensions of an array defined

mathematical expressions obtained by mathematical operator $\bar{\&}$



Types of sub-arrays (2)

`int c[2][3][4];`

<code>(c+i)</code>	<code>= &c[i]</code>
<code>((c+i)+j)</code>	<code>= &c[i]</code>
<code>((c+i)+j)+k)</code>	<code>= &c[i]</code>

<code>v(c+i)</code>	<code>= v(c) + i * sizeof(3*4)</code>
<code>v((c+i)+j)</code>	<code>= v(c+i) + j * sizeof(4)</code>
<code>v(((c+i)+j)+k)</code>	<code>= v((c+i)+j) + k</code>

`int c[2][3][4];`

<code>&c[i]</code>	<code>:: int (*) [3][4]</code>
<code>((c+i)+j)</code>	<code>:: int (*) [4]</code>
<code>((c+i)+j)+k)</code>	<code>:: int (*)</code>

<code>v(c+i)</code>	<code>= v(c) + i * sizeof(3*4)</code>
<code>v((c+i)+j)</code>	<code>= v(c+i) + j * sizeof(4)</code>
<code>v(((c+i)+j)+k)</code>	<code>= v((c+i)+j) + k</code>

Types of sub-arrays (2)

`int c[2][3][4];`

<code>(c+i)</code>	<code>= &c[i]</code>
<code>((c+i)+j)</code>	<code>= &c[i]</code>
<code>((c+i)+j)+k)</code>	<code>= &c[i]</code>

<code>v(c+i)</code>	<code>= v(c) + i * sizeof(3*4)</code>
<code>v((c+i)+j)</code>	<code>= v(c+i) + j * sizeof(4)</code>
<code>v(((c+i)+j)+k)</code>	<code>= v((c+i)+j) + k</code>

`int c[2][3][4];`

<code>&c[i]</code>	<code>:: int (*) [3][4]</code>
<code>((c+i)+j)</code>	<code>:: int (*) [4]</code>
<code>((c+i)+j)+k)</code>	<code>:: int (*)</code>

<code>v(c+i)</code>	<code>= v(c) + i * sizeof(3*4)</code>
<code>v((c+i)+j)</code>	<code>= v(c+i) + j * sizeof(4)</code>
<code>v(((c+i)+j)+k)</code>	<code>= v((c+i)+j) + k</code>

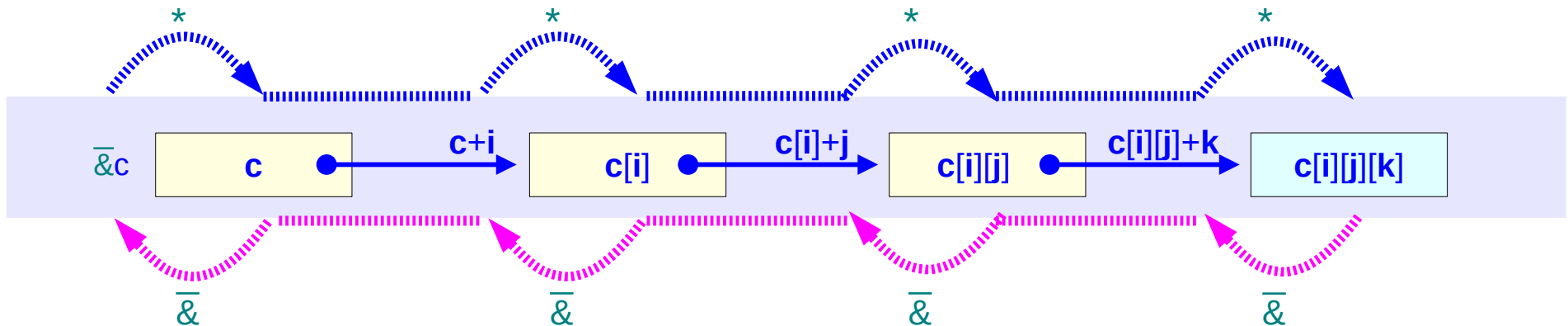
Recursive applications of $*$ and $\bar{\&}$ are possible

$$*(*(*(\mathbf{c}+\mathbf{i})+\mathbf{j})+\mathbf{k})]$$

$$\begin{aligned} & *(*(*(\mathbf{c}+\mathbf{i})+\mathbf{j})+\mathbf{k}) = \\ & *(*(\mathbf{c}[\mathbf{i}]+\mathbf{j})+\mathbf{k}) = \\ & *(\mathbf{c}[\mathbf{i}][\mathbf{j}]+\mathbf{k}) = \\ & \mathbf{c}[\mathbf{i}][\mathbf{j}][\mathbf{k}] \end{aligned}$$

$$\bar{\&}(\bar{\&}(\bar{\&}(\mathbf{c}[\mathbf{i}])[\mathbf{j}])[\mathbf{k}])]$$

$$\begin{aligned} & \bar{\&}(\bar{\&}(\bar{\&}(\mathbf{c}[\mathbf{i}])[\mathbf{j}])[\mathbf{k}]) = \\ & (\bar{\&}(\bar{\&}(\mathbf{c}[\mathbf{i}])[\mathbf{j}])+\mathbf{k}) = \\ & ((\bar{\&}(\mathbf{c}[\mathbf{i}])+\mathbf{j})+\mathbf{k}) = \\ & (((\mathbf{c}+\mathbf{i})+\mathbf{j})+\mathbf{k}) \end{aligned}$$



Outward and Inward Recursive applications of $*$ and $\bar{\&}$

←

$$*(*(*(\mathbf{c}+\mathbf{i})+\mathbf{j})+\mathbf{k})]$$

↓

$$\begin{aligned} &*(*(*(\mathbf{c}+\mathbf{i})+\mathbf{j})+\mathbf{k}) = \\ &*(*(\mathbf{c}[\mathbf{i}]+\mathbf{j})+\mathbf{k}) = \\ &* (\mathbf{c}[\mathbf{i}][\mathbf{j}]+\mathbf{k}) = \\ &\mathbf{c}[\mathbf{i}][\mathbf{j}][\mathbf{k}] \end{aligned}$$

→

$$*(*(*(\mathbf{c}+\mathbf{i})+\mathbf{j})+\mathbf{k})]$$

↓

$$\begin{aligned} &*(*(*(\mathbf{c}+\mathbf{i})+\mathbf{j})+\mathbf{k}) = \\ &(*(*(\mathbf{c}+\mathbf{i})+\mathbf{j})[\mathbf{k}]) = \\ &(*(\mathbf{c}+\mathbf{i}))[\mathbf{j}][\mathbf{k}] = \\ &\mathbf{c}[\mathbf{i}][\mathbf{j}][\mathbf{k}] \end{aligned}$$

←

$$\bar{\&}(\bar{\&}(\bar{\&}(\mathbf{c}[\mathbf{i}])[\mathbf{j}]))[\mathbf{k}]$$

↓

$$\begin{aligned} &\bar{\&}(\bar{\&}(\bar{\&}(\mathbf{c}[\mathbf{i}])[\mathbf{j}]))[\mathbf{k}] = \\ &\bar{\&}(\bar{\&}((\mathbf{c}+\mathbf{i})[\mathbf{j}]))[\mathbf{k}] = \\ &\bar{\&}(((\mathbf{c}+\mathbf{i})+\mathbf{j})[\mathbf{k}]) = \\ &(((\mathbf{c}+\mathbf{i})+\mathbf{j})+\mathbf{k}) \end{aligned}$$

→

$$\bar{\&}(\bar{\&}(\bar{\&}(\mathbf{c}[\mathbf{i}])[\mathbf{j}]))[\mathbf{k}]$$

↓

$$\begin{aligned} &\bar{\&}(\bar{\&}(\bar{\&}(\mathbf{c}[\mathbf{i}])[\mathbf{j}]))[\mathbf{k}] = \\ &(\bar{\&}(\bar{\&}(\mathbf{c}[\mathbf{i}])[\mathbf{j}])+\mathbf{k}) = \\ &((\bar{\&}(\mathbf{c}[\mathbf{i}])+\mathbf{j})+\mathbf{k}) = \\ &(((\mathbf{c}+\mathbf{i})+\mathbf{j})+\mathbf{k}) \end{aligned}$$

mathematical expressions

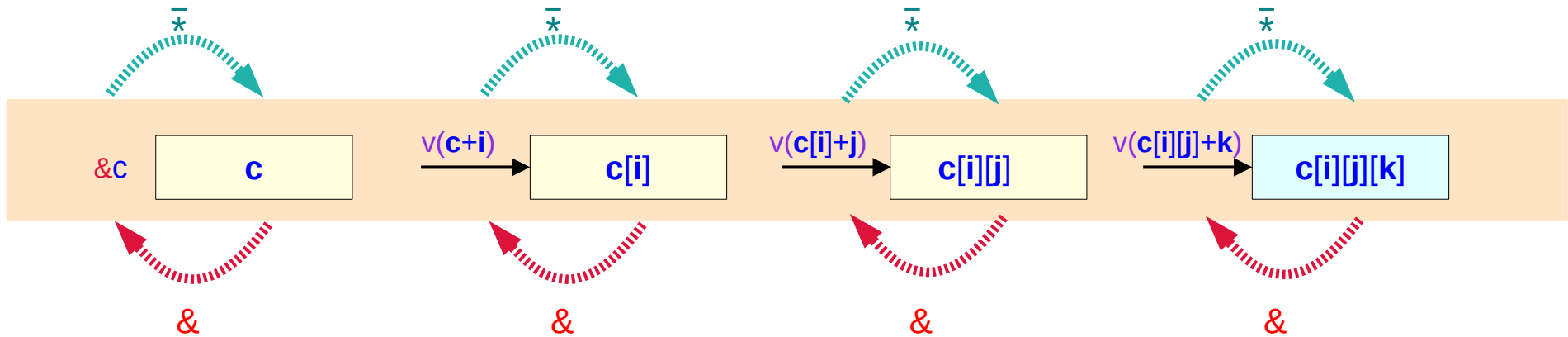
Recursive applications of $\bar{*}$ and $\&$ are not possible (2)

~~$\bar{*}(\bar{*}(\bar{*}(c+i)+j)+k)$~~

not allowed
even in mathematical expressions

~~$\&(\&(\&(c[i])[j])[k])$~~

not allowed
in C expressions



Recursive applications of $*$ and $\overline{\&}$

legal C exp

$*(*(*(\mathbf{c}+\mathbf{i})+\mathbf{j})+\mathbf{k})$

legal math. exp

$(((\mathbf{c}+\mathbf{i})+\mathbf{j})+\mathbf{k})$

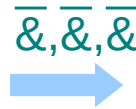


legal C exp

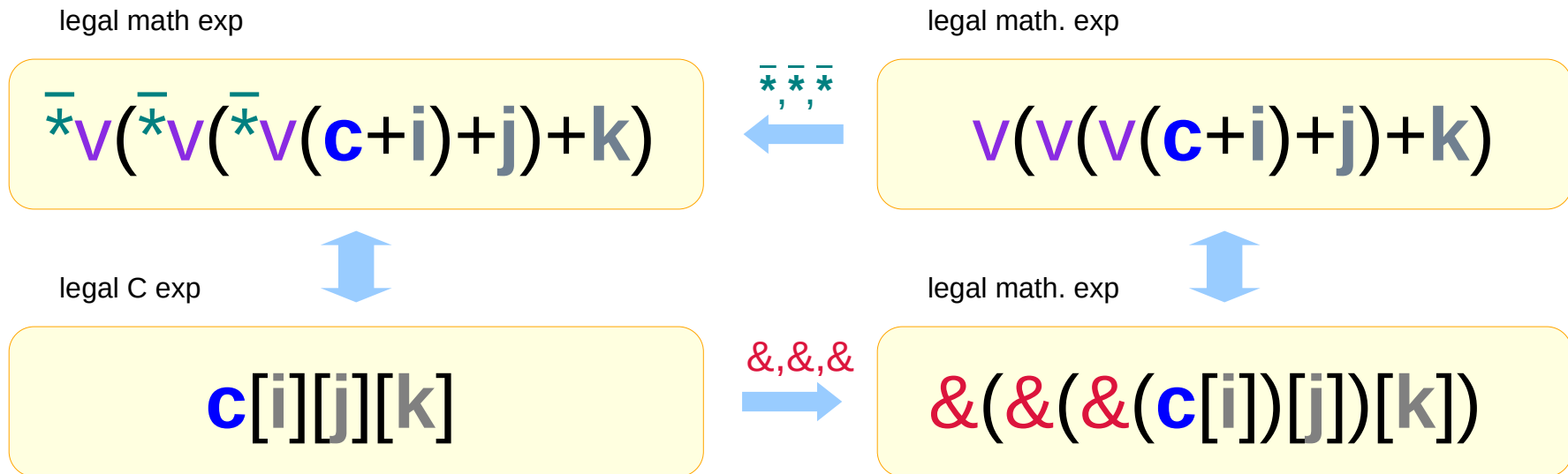
$\mathbf{c}[\mathbf{i}][\mathbf{j}][\mathbf{k}]$

legal math. exp

$\overline{\&}(\overline{\&}(\overline{\&}(\mathbf{c}[\mathbf{i}])[\mathbf{j}])[\mathbf{k}])$



Recursive applications of $\bar{*}$ and $\&$



Recursive applications of $\overline{*}$ and $\&$ are not possible (3)

legal C expression

$$*(*(*(\mathbf{c}+\mathbf{i})+\mathbf{j})+\mathbf{k})$$

legal math. expression

$$\overline{*}v(\overline{*}v(\overline{*}v(\mathbf{c}+\mathbf{i})+\mathbf{j})+\mathbf{k})$$

illegal math. expression

$$\overline{*}(\overline{*}(\overline{*}(\mathbf{c}+\mathbf{i})+\mathbf{j})+\mathbf{k})$$

legal math. expression

$$\overline{\&}(\overline{\&}(\overline{\&}(\mathbf{c}[\mathbf{i}])[\mathbf{j}])[\mathbf{k}])$$

illegal math. expression

$$v\overline{\&}(v\overline{\&}(v\overline{\&}(\mathbf{c}[\mathbf{i}])[\mathbf{j}])[\mathbf{k}])$$

illegal C expression

$$\&(\&(\&(\mathbf{c}[\mathbf{i}])[\mathbf{j}])[\mathbf{k}])$$

Recursive applications of $\overline{*}$ and $\&$ are not possible (3)

legal C expression

$$*(*(*(\mathbf{c}+\mathbf{i})+\mathbf{j})+\mathbf{k})]$$

legal math. expression

$$\overline{\&}(\overline{\&}(\overline{\&}(\mathbf{c}[\mathbf{i}])[\mathbf{j}])[\mathbf{k}])$$

legal math. expression

$$\overline{*}\mathbf{v}(\overline{*}\mathbf{v}(\overline{*}\mathbf{v}(\mathbf{c}+\mathbf{i})+\mathbf{j})+\mathbf{k})$$

legal math. expression

$$\overline{*}\&(\overline{*}\&(\overline{*}\&(\mathbf{c}[\mathbf{i}])[\mathbf{j}])[\mathbf{k}])$$

Recursive applications of $\bar{*}$ and $\&$ are not possible (2)

$$\begin{aligned} c[i][j][k] &= \bar{*}v(c[i][j] + k) \\ &= \bar{*}v(\bar{*}v(c[i] + j) + k) \\ &= \bar{*}v(\bar{*}v(\bar{*}v(c + i) + j) + k) \end{aligned}$$

$$\begin{aligned} v(v(v(c + i) + j) + k) &= (\&(c[i]) + j) + k \\ &= (\&(\&(c[i])[j])) + k \\ &= \&(\&(\&(c[i])[j])[k]) \end{aligned}$$

$$\begin{aligned} a[i] &= \&(c[i]) \\ b[i][j] &= (\&(\&(c[i])[j])) \end{aligned}$$

legal math. expression

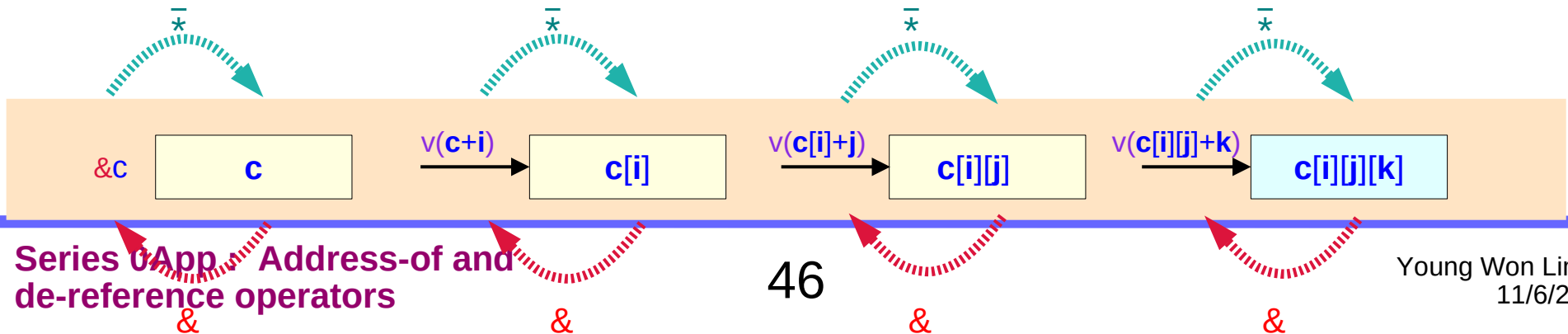
$$\bar{\&}(\bar{\&}(\bar{\&}(c[i])[j])[k])$$

legal math. expression

$$\bar{*}v\bar{\&}(\bar{*}v\bar{\&}(\bar{*}v\bar{\&}(c[i])[j])[k])$$

illegal C expression

$$\bar{*}\&(\bar{*}\&(\bar{*}\&(c[i])[j])[k])$$



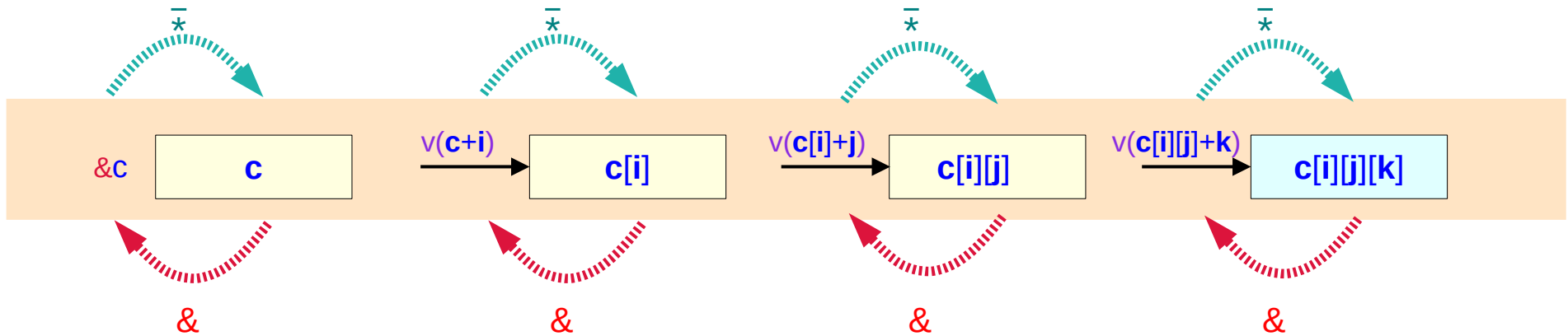
Recursive applications of $\bar{*}$ and $\&$ are not possible (1)

$\bar{*}v(\bar{*}v(\bar{*}v(c+i)+j)+k)$

$\begin{aligned}\bar{*}v(\bar{*}v(\bar{*}v(c+i)+j)+k) &= \\ \bar{*}v(\bar{*}v(c[i]+j)+k) &= \\ \bar{*}v(c[i][j]+k) &= \\ c[i][j][k] &\end{aligned}$

~~$v\bar{\&}(v\bar{\&}(v\bar{\&}(c[i])[j])[k])$~~

~~$\&(\&(\&(c[i])[j])[k])$~~



Recursive applications of $\bar{*}$ and $\&$ are not possible (2)

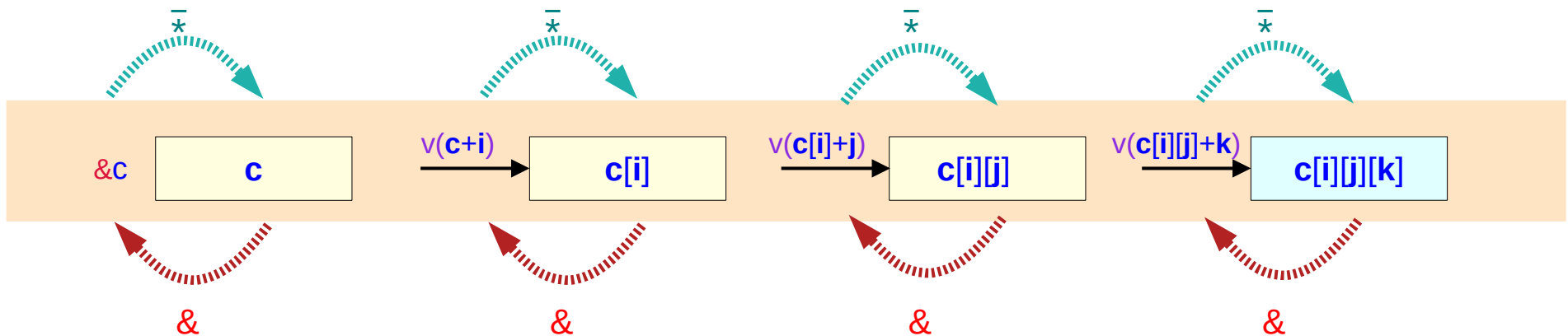
~~$\&(\&(\&(\mathbf{c}[i])[j])[k])$~~

~~$\mathbf{v}\bar{\&}(\mathbf{v}\bar{\&}(\mathbf{v}\bar{\&}(\mathbf{c}[i])[j])[k])$~~

$$\begin{aligned}\&(\bar{\&}(\bar{\&}(\mathbf{c}[i])[j])[k]) &= \\ \mathbf{v}\bar{\&}(\bar{\&}(\bar{\&}(\mathbf{c}[i])[j])[k]) &= \\ \mathbf{v}(\bar{\&}(\bar{\&}(\mathbf{c}[i])[j])+k) &= \\ \mathbf{v}((\bar{\&}(\mathbf{c}[i])+j)+k) &= \\ \mathbf{v}(((\mathbf{c}+i)+j)+k) &= \end{aligned}$$

$$\begin{aligned}\&(\bar{\&}(\mathbf{c}[i])[j])[k] &= \\ \mathbf{v}\bar{\&}(\bar{\&}(\mathbf{c}[i])[j])[k] &= \\ \mathbf{v}((\bar{\&}(\mathbf{c}[i])+j)[k] &= \\ \mathbf{v}((\mathbf{c}+i)+j)[k] &= \end{aligned}$$

$$\begin{aligned}\&(\mathbf{c}[i])[j][k] &= \\ \mathbf{v}\bar{\&}(\mathbf{c}[i])[j][k] &= \\ \mathbf{v}((\mathbf{c}+i)[j])[k] &= \end{aligned}$$



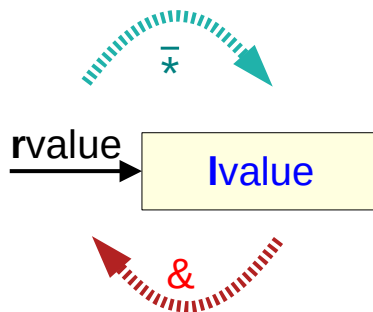
Recursive applications of $\overline{*}$ and $\&$ are not possible (3)

~~$\&(\&(\&(c[i])[j])[k])$~~

$\&$ takes only **lvalue** variable
returns **rvalue** address value

successive application of $\&$ is not possible,
unless intermediate variables and proper
type casting

$\overline{*}$ must be applied to an **rvalue** address
value, $\overline{*}$ operator cannot be
applied successively.

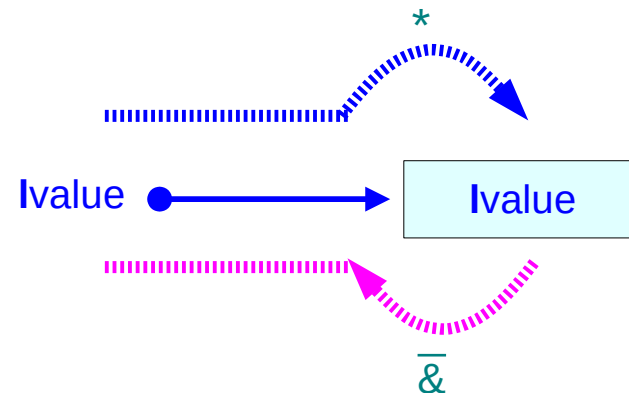


$\overline{\&}(\overline{\&}(\overline{\&}(c[i])[j])[k])$

$\overline{\&}$ takes a dereferenced **lvalue** variable
returns **lvalue** variable

successive application of $\overline{\&}$ is possible

but, $*p$ becomes a **lvalue** variable
 $*$ operator can be applied successively.



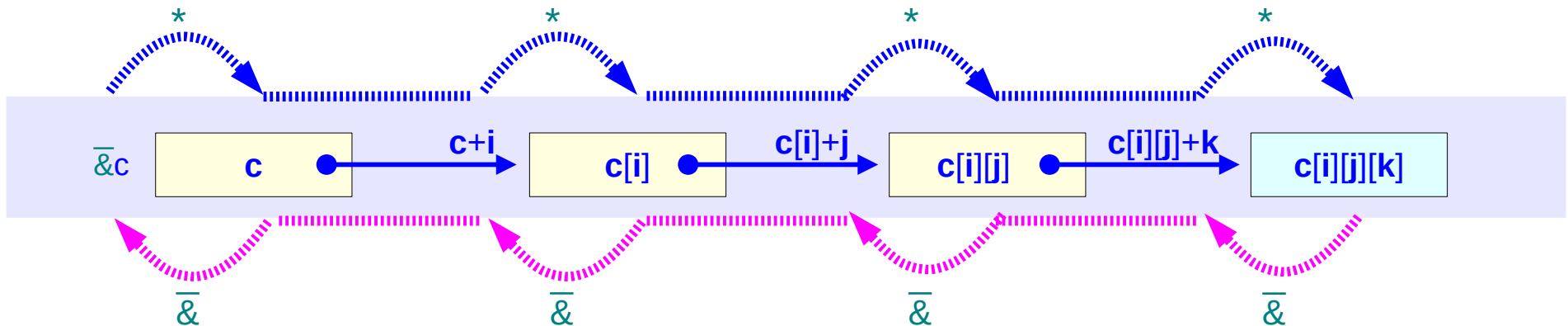
Inverse Relations between $*$ and $\overline{\&}$

$$\begin{aligned} \mathbf{c[i][j][k]} &= *(\mathbf{c[i][j]} + \mathbf{k}) \\ \mathbf{c[i][j]} &= *(\mathbf{c[i]} + \mathbf{j}) \\ \mathbf{c[i]} &= *(\mathbf{c} + \mathbf{i}) \end{aligned}$$

$$\begin{aligned} \mathbf{c[i][j][k]} &= *(\mathbf{c[i][j]} + \mathbf{k}) \\ &= *(*(\mathbf{c[i]} + \mathbf{j}) + \mathbf{k}) \\ &= *(*(*(\mathbf{c} + \mathbf{i}) + \mathbf{j}) + \mathbf{k}) \end{aligned}$$

$$\begin{aligned} \overline{\&}(\mathbf{c[i][j][k]}) &= \mathbf{c[i][j]} + \mathbf{k} \\ \overline{\&}(\mathbf{c[i][j]}) &= \mathbf{c[i]} + \mathbf{j} \\ \overline{\&}(\mathbf{c[i]}) &= \mathbf{c} + \mathbf{i} \end{aligned}$$

$$\begin{aligned} ((\mathbf{c} + \mathbf{i}) + \mathbf{j}) + \mathbf{k} &= (\overline{\&}(\mathbf{c[i]}) + \mathbf{j}) + \mathbf{k} \\ &= (\overline{\&}(\overline{\&}(\mathbf{c[i]})[\mathbf{j}])) + \mathbf{k} \\ &= \overline{\&}(\overline{\&}(\overline{\&}(\mathbf{c[i]})[\mathbf{j}])[\mathbf{k}]) \end{aligned}$$



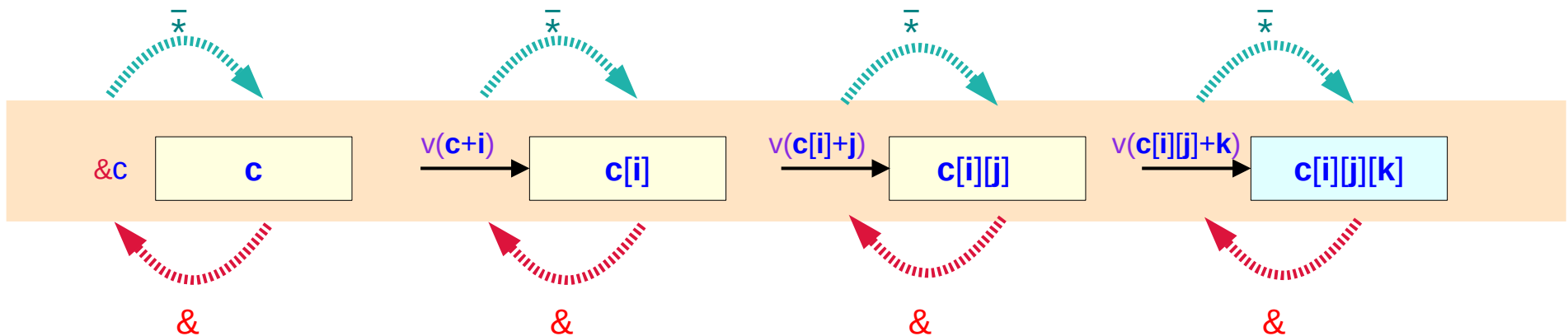
Inverse Relations between $\bar{*}$ and $\&$

$$\begin{aligned} c[i][j][k] &= \bar{*}v(c[i][j] + k) \\ c[i][j] &= \bar{*}v(c[i] + j) \\ c[i] &= \bar{*}v(c + i) \end{aligned}$$

$$\begin{aligned} c[i][j][k] &= \bar{*}v(c[i][j] + k) \\ &= \bar{*}v(\bar{*}v(c[i] + j) + k) \\ &= \bar{*}v(\bar{*}v(\bar{*}v(c + i) + j) + k) \end{aligned}$$

$$\begin{aligned} \&(c[i][j][k]) &= v(c[i][j] + k) \\ \&(c[i][j]) &= v(c[i] + j) \\ \&(c[i]) &= v(c + i) \end{aligned}$$

$$\begin{aligned} v(v(v(c + i) + j) + k) &= (\&(c[i]) + j) + k \\ &= (\&(\&(c[i])[j])) + k \\ &= \&(\&(\&(c[i])[j])[k]) \end{aligned}$$



Inverse Relations between $\bar{*}$ and $\&$

$$\begin{aligned} \mathbf{c[i][j][k]} &= \bar{*}\mathbf{v}(\mathbf{c[i][j]} + \mathbf{k}) \\ \mathbf{c[i][j]} &= \bar{*}\mathbf{v}(\mathbf{c[i]} + \mathbf{j}) \\ \mathbf{c[i]} &= \bar{*}\mathbf{v}(\mathbf{c} + \mathbf{i}) \end{aligned}$$

$$\begin{aligned} \mathbf{c[i][j][k]} &= \bar{*}\mathbf{v}(\mathbf{c[i][j]} + \mathbf{k}) \\ &= \bar{*}\mathbf{v}(\bar{*}\mathbf{v}(\mathbf{c[i]} + \mathbf{j}) + \mathbf{k}) \\ &= \bar{*}\mathbf{v}(\bar{*}\mathbf{v}(\bar{*}\mathbf{v}(\mathbf{c} + \mathbf{i}) + \mathbf{j}) + \mathbf{k}) \end{aligned}$$

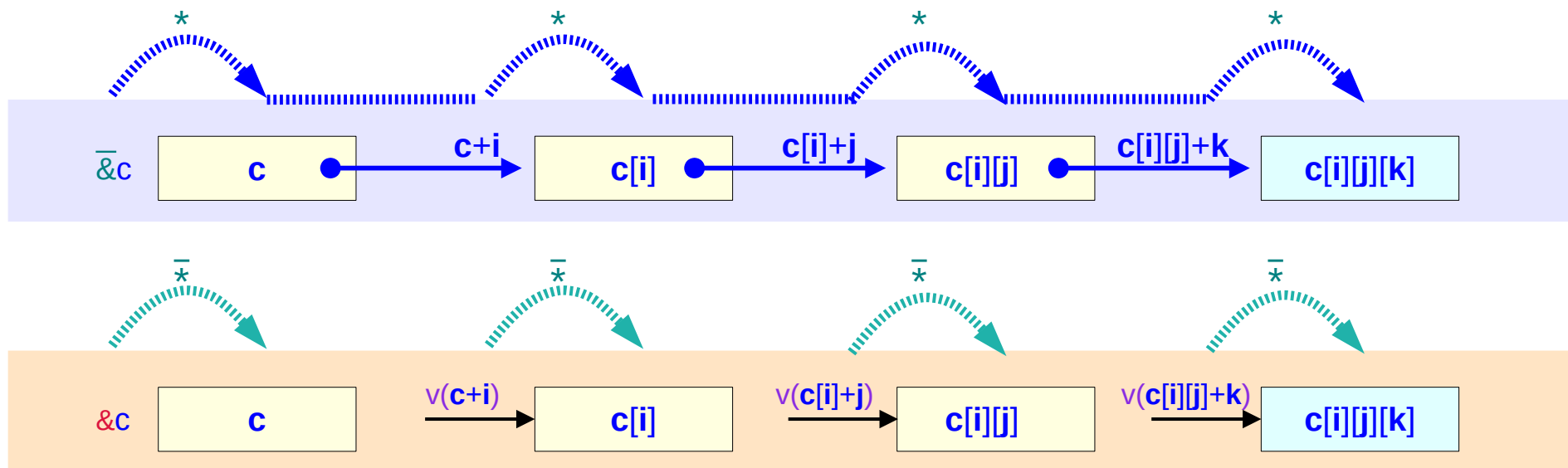
$$\begin{aligned} \&(\mathbf{c[i][j][k]}) &= \mathbf{v}(\mathbf{c[i][j]} + \mathbf{k}) \\ \&(\mathbf{c[i][j]}) &= \mathbf{v}(\mathbf{c[i]} + \mathbf{j}) \\ \&(\mathbf{c[i]}) &= \mathbf{v}(\mathbf{c} + \mathbf{i}) \end{aligned}$$

$$\begin{aligned} \mathbf{v}(\mathbf{v}(\mathbf{v}(\mathbf{c} + \mathbf{i}) + \mathbf{j}) + \mathbf{k}) &= (\&(\mathbf{c[i]}) + \mathbf{j}) + \mathbf{k} \\ &= (\&(\&(\mathbf{c[i]})[\mathbf{j}])) + \mathbf{k} \\ &= \&(\&(\&(\mathbf{c[i]})[\mathbf{j}])[\mathbf{k}]) \end{aligned}$$

$$\begin{aligned} \mathbf{v}(\mathbf{v}(\mathbf{v}(\mathbf{c} + \mathbf{i}) + \mathbf{j}) + \mathbf{k}) &= \&(\mathbf{v}(\mathbf{v}(\mathbf{c} + \mathbf{i}) + \mathbf{j})[\mathbf{k}]) \\ &= \&(\&(\mathbf{v}(\mathbf{c} + \mathbf{i})[\mathbf{j}])[\mathbf{k}]) \\ &= \&(\&(\&(\mathbf{c[i]})[\mathbf{j}])[\mathbf{k}]) \end{aligned}$$

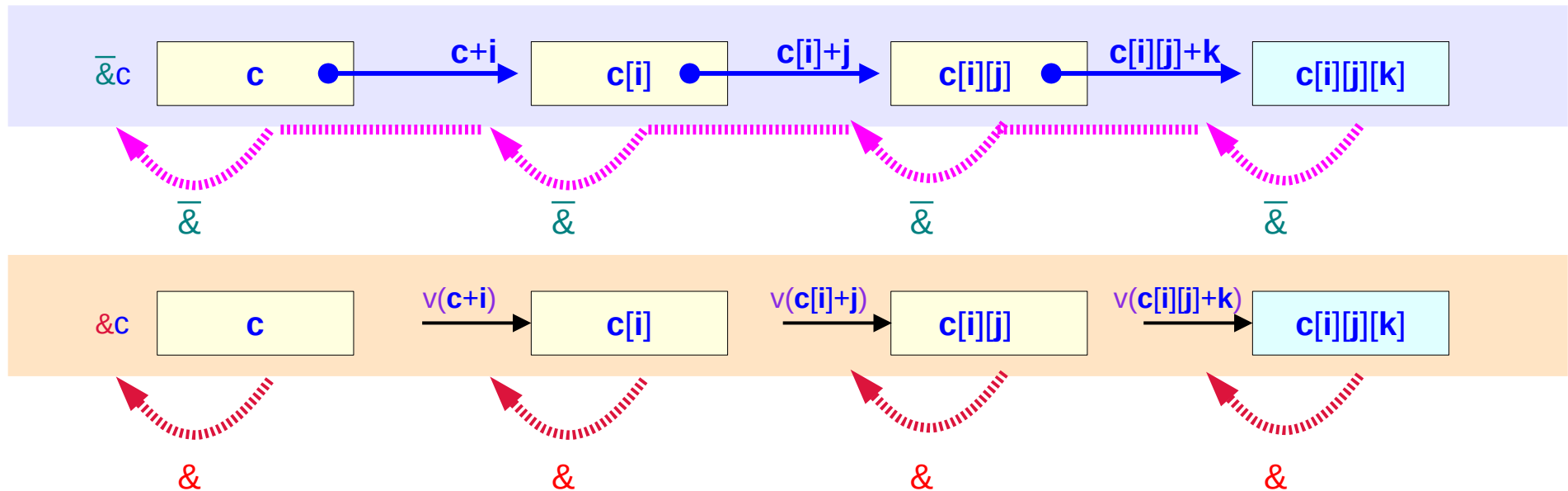
$$\begin{aligned}
 \mathbf{c[i][j][k]} &= \mathbf{*}(\mathbf{c[i][j]} + \mathbf{k}) \\
 \mathbf{c[i][j]} &= \mathbf{*}(\mathbf{c[i]} + \mathbf{j}) \\
 \mathbf{c[i]} &= \mathbf{*}(\mathbf{c} + \mathbf{i})
 \end{aligned}$$

$$\begin{aligned}
 \mathbf{c[i][j][k]} &= \mathbf{\bar{*}v(c[i][j]} + \mathbf{k}) &= \mathbf{\bar{*}(value(c[i][j]) + k * sizeof(c[0][0][0]))} \\
 \mathbf{c[i][j]} &= \mathbf{\bar{*}v(c[i]} + \mathbf{j}) &= \mathbf{\bar{*}(value(c[i]) + j * sizeof(c[0][0]))} \\
 \mathbf{c[i]} &= \mathbf{\bar{*}v(c} + \mathbf{i}) &= \mathbf{\bar{*}(value(c) + i * sizeof(c[0]))}
 \end{aligned}$$



$$\begin{aligned}\overline{\&(\mathbf{c}[\mathbf{i}][\mathbf{j}][\mathbf{k}])} &= \mathbf{c}[\mathbf{i}][\mathbf{j}] + \mathbf{k} \\ \overline{\&(\mathbf{c}[\mathbf{i}][\mathbf{j}])} &= \mathbf{c}[\mathbf{i}] + \mathbf{j} \\ \overline{\&(\mathbf{c}[\mathbf{i}])} &= \mathbf{c} + \mathbf{i}\end{aligned}$$

$$\begin{aligned}\&(\mathbf{c}[\mathbf{i}][\mathbf{j}][\mathbf{k}]) &= \mathbf{v}(\mathbf{c}[\mathbf{i}][\mathbf{j}] + \mathbf{k}) &= (\text{value}(\mathbf{c}[\mathbf{i}][\mathbf{j}]) + \mathbf{k} * \text{sizeof}(\mathbf{c}[0][0][0])) \\ \&(\mathbf{c}[\mathbf{i}][\mathbf{j}]) &= \mathbf{v}(\mathbf{c}[\mathbf{i}] + \mathbf{j}) &= (\text{value}(\mathbf{c}[\mathbf{i}]) + \mathbf{j} * \text{sizeof}(\mathbf{c}[0][0])) \\ \&(\mathbf{c}[\mathbf{i}]) &= \mathbf{v}(\mathbf{c} + \mathbf{i}) &= (\text{value}(\mathbf{c}) + \mathbf{i} * \text{sizeof}(\mathbf{c}[0]))\end{aligned}$$



Two step deferencing in type II (1) – without skipping

$$\begin{aligned}
 \overline{\&}(c[i][j][k]) &= \overline{\&}(* (c[i][j] + k)) &= c[i][j] &+ k \\
 \overline{\&}(c[i][j]) &= \overline{\&}(* (c[i] + j)) &= c[i] &+ j \\
 \overline{\&}(c[i]) &= \overline{\&}(* (c + i)) &= c &+ i
 \end{aligned}$$

$$\begin{aligned}
 *\overline{\&}(c[i][j][k]) &= *\overline{\&}(* (c[i][j] + k)) &= *(c[i][j] + k) \\
 *\overline{\&}(c[i][j]) &= *\overline{\&}(* (c[i] + j)) &= *(c[i] + j) \\
 *\overline{\&}(c[i]) &= *\overline{\&}(* (c + i)) &= *(c + i)
 \end{aligned}$$

$$\begin{aligned}
 \&(c[i][j][k]) &= \&(* (c[i][j] + k)) &= c[i][j] &+ k * \text{sizeof}(c[0][0][0]) \\
 \&(c[i][j]) &= \&(* (c[i] + j)) &= c[i] &+ j * \text{sizeof}(c[0][0]) \\
 \&(c[i]) &= \&(* (c + i)) &= c &+ i * \text{sizeof}(c[0])
 \end{aligned}$$

$$\begin{aligned}
 c[i][j][k] &= *\overline{\&}(* (c[i][j] + k)) &= \overline{\&}(c[i][j] + k * \text{sizeof}(c[0][0][0])) \\
 c[i][j] &= *\overline{\&}(* (c[i] + j)) &= \overline{\&}(c[i] + j * \text{sizeof}(c[0][0])) \\
 c[i] &= *\overline{\&}(* (c + i)) &= \overline{\&}(c + i * \text{sizeof}(c[0]))
 \end{aligned}$$

Two step deferencing in type II (1) – without skipping

$$\begin{aligned}\overline{\&}(c[i][j][k]) &= c[i][j] + k \\ \overline{\&}(c[i][j]) &= c[i] + j \\ \overline{\&}(c[i]) &= c + i\end{aligned}$$

$$\begin{aligned}c[i][j][k] &= *(c[i][j] + k) \\ c[i][j] &= *(c[i] + j) \\ c[i] &= *(c + i)\end{aligned}$$

$$\begin{aligned}\&(c[i][j][k]) &= c[i][j] + k * \text{sizeof}(c[0][0][0]) \\ \&(c[i][j]) &= c[i] + j * \text{sizeof}(c[0][0]) \\ \&(c[i]) &= c + i * \text{sizeof}(c[0])\end{aligned}$$

$$\begin{aligned}c[i][j][k] &= \overline{*}(c[i][j] + k * \text{sizeof}(c[0][0][0])) \\ c[i][j] &= \overline{*}(c[i] + j * \text{sizeof}(c[0][0])) \\ c[i] &= \overline{*}(c + i * \text{sizeof}(c[0]))\end{aligned}$$

Two step deferencing in type II (1) – without skipping

$$\begin{aligned}\overline{\&}(c[i][j][k]) &= c[i][j] + k \\ \overline{\&}(c[i][j]) &= c[i] + j \\ \overline{\&}(c[i]) &= c + i\end{aligned}$$

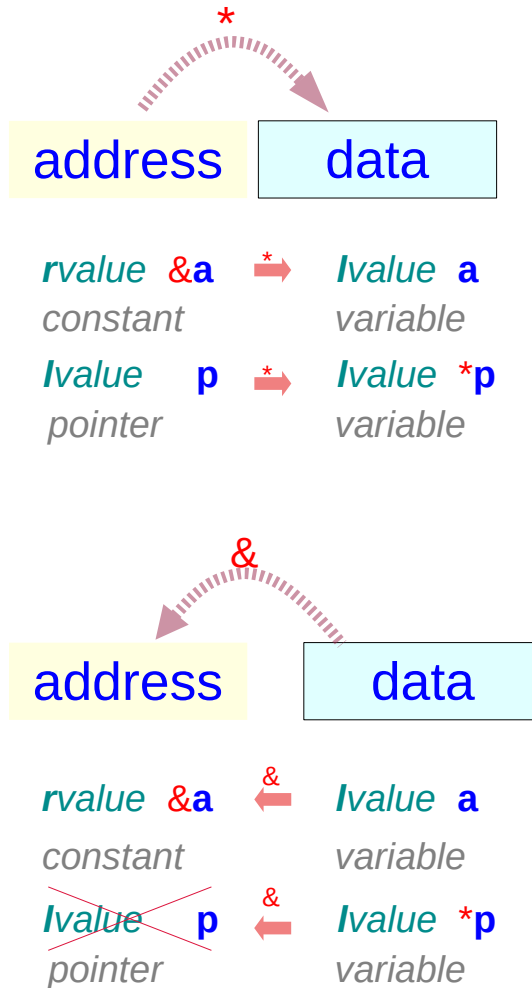
$$\begin{aligned}\&(c[i][j][k]) &= \text{value}(c[i][j] + k) \\ \&(c[i][j]) &= \text{value}(c[i] + j) \\ \&(c[i]) &= \text{value}(c + i)\end{aligned}$$

$$\begin{aligned}c[i][j][k] &= *(c[i][j] + k) \\ c[i][j] &= *(c[i] + j) \\ c[i] &= *(c + i)\end{aligned}$$

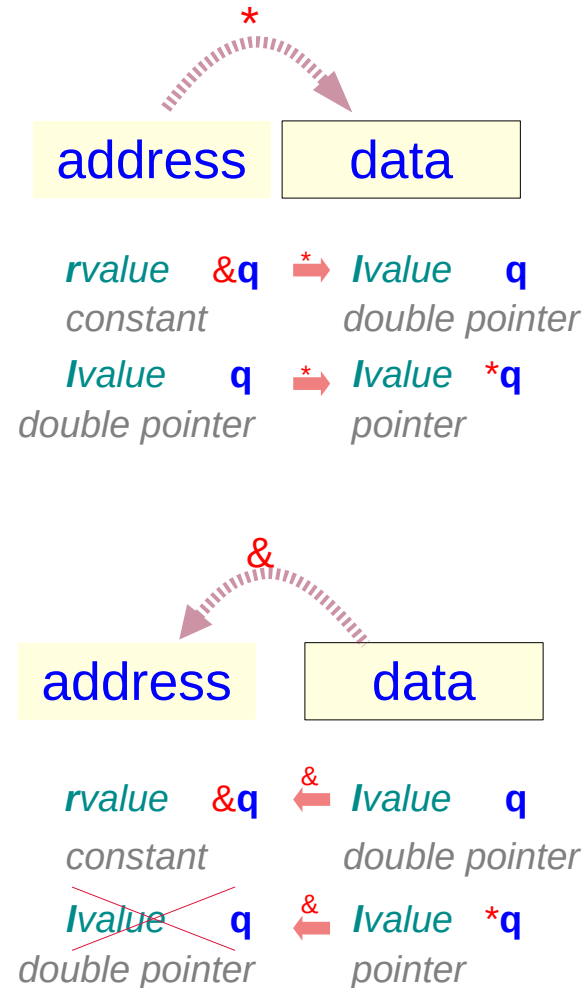
$$\begin{aligned}c[i][j][k] &= \overline{*} \text{value}(c[i][j] + k) \\ c[i][j] &= \overline{*} \text{value}(c[i] + j) \\ c[i] &= \overline{*} \text{value}(c + i)\end{aligned}$$

Address-of & and dereference * C operators (1)

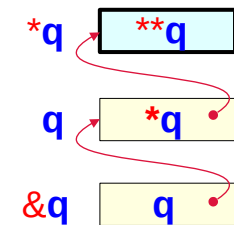
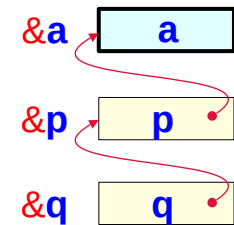
Primitive Data Type



Pointer Data Type

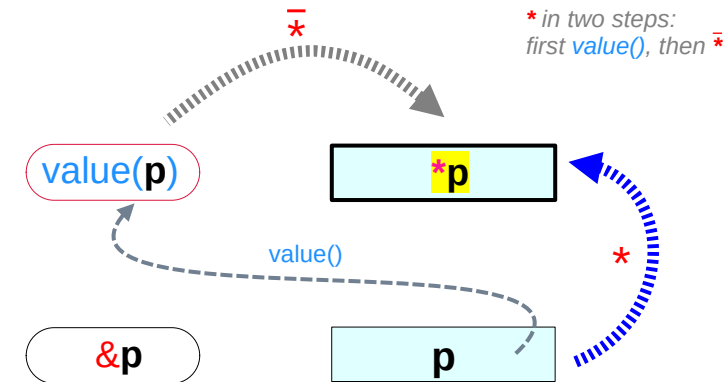
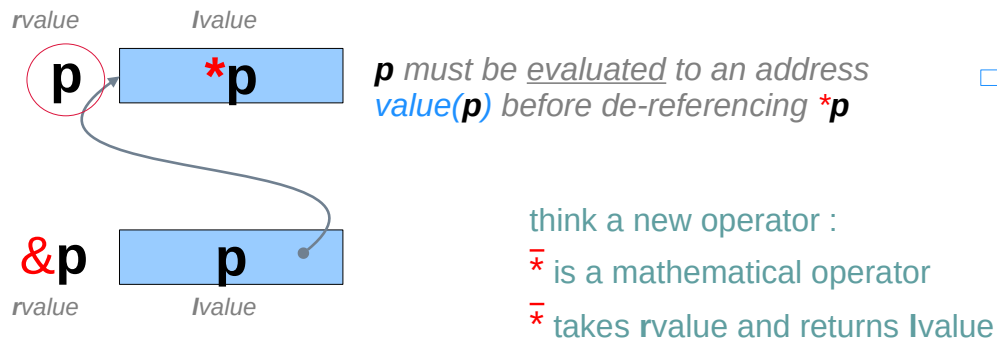


```
int a ;  
int * p ;  
int ** q ;
```

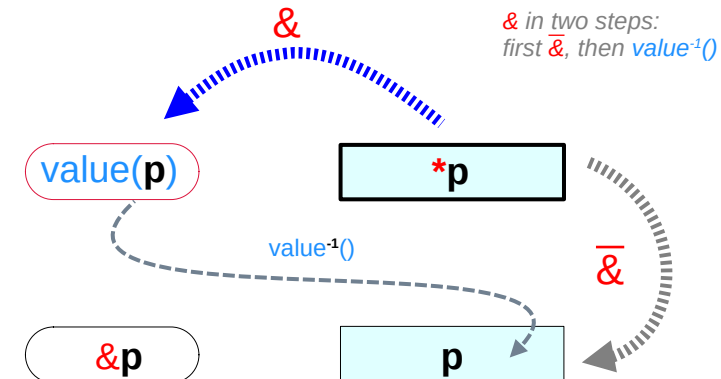
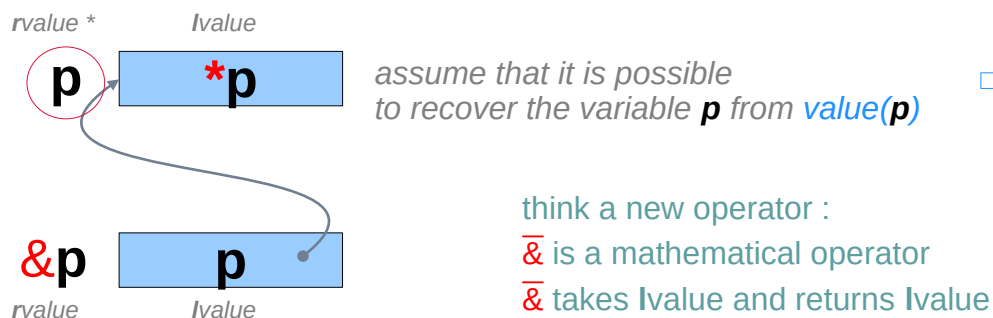


& and * operators in pointer de-referencing

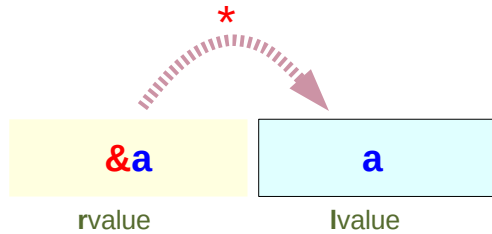
Two step De-reference * operation



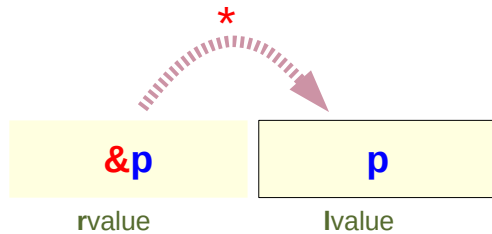
Two step Address-of & operation



Type, Size, and Value attributes of an **lvalue**



lvalue ← **rvalue*
a ***&a**



lvalue ← **rvalue*
p ***&p**

lvalue is associated with a memory location

lvalue has the following attributes

- Type
- Size
- Value

rvalue has the only attribute

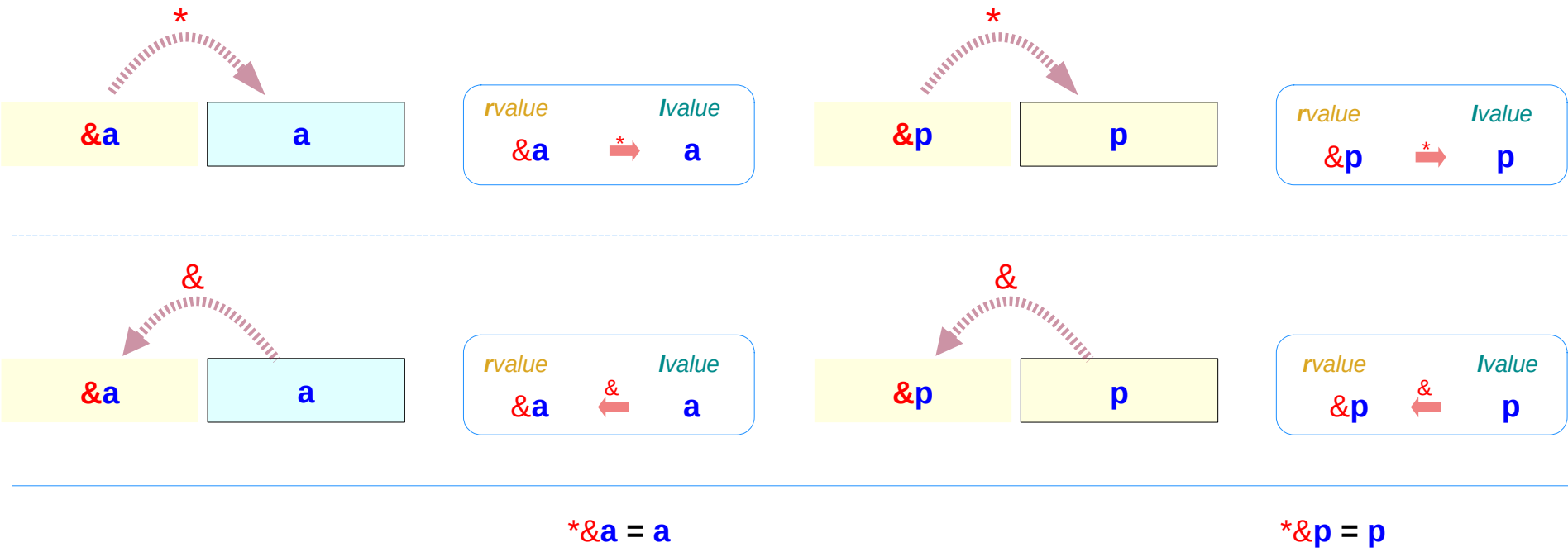
- Value

assume the function `value()`

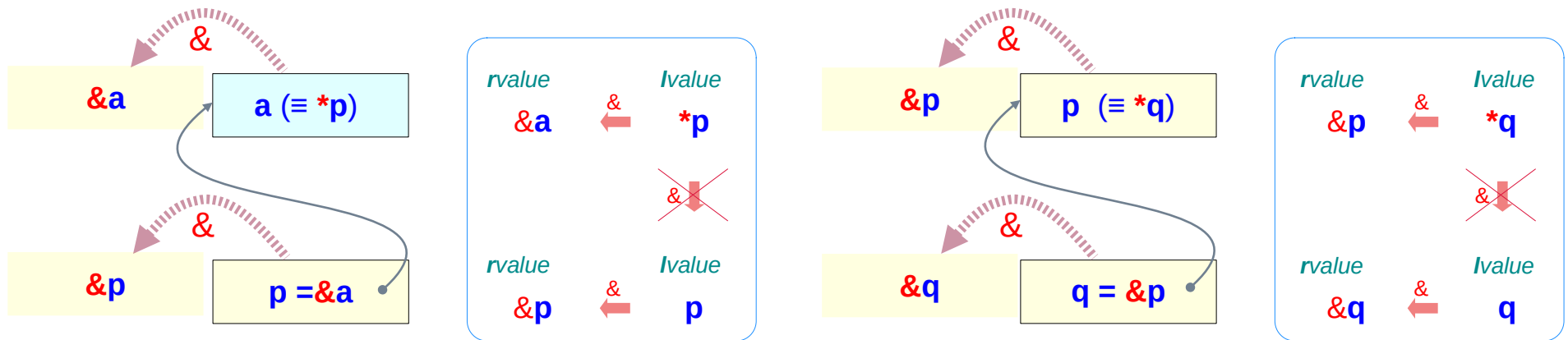
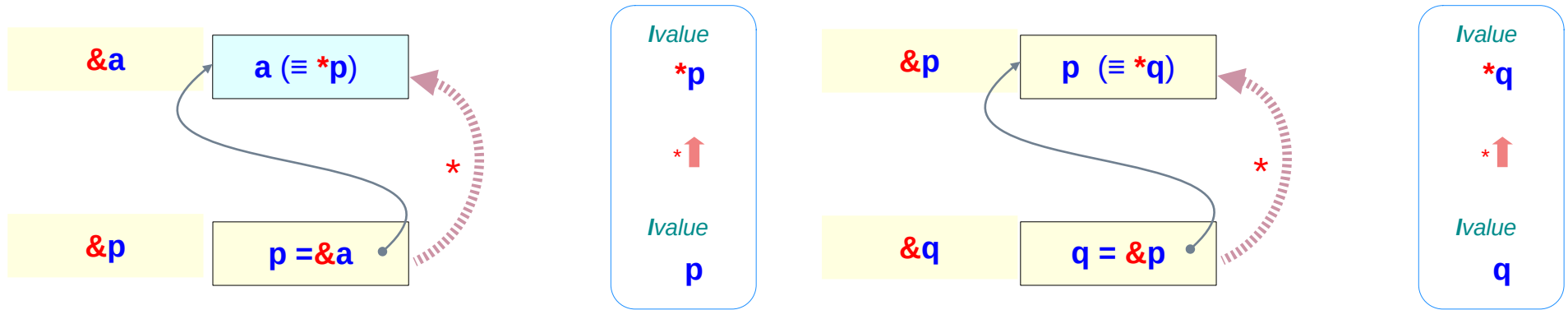
`value(lvalue)` returns
the **Value** attribute of **lvalue**

`value(rvalue)` returns
the **Value** attribute of **rvalue**

Address-of & and dereference * C operators (1)



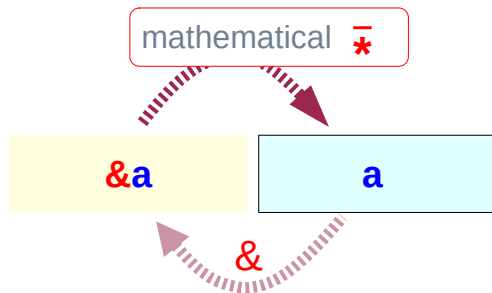
Address-of & and dereference * C operators (2)



`*&p = p`
~~`&*p = p`~~ *value(p)*

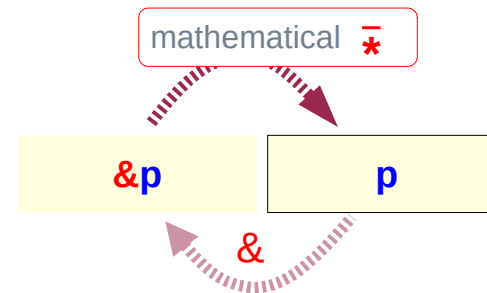
`*&q = q`
~~`&*q = q`~~ *value(q)*

Introducing mathematical operators : $\bar{\&}$ and $\bar{*}$



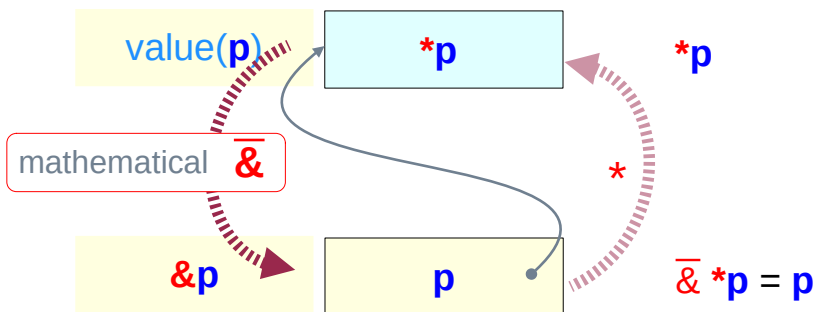
$$\bar{*} \&a = a$$

$\&a$

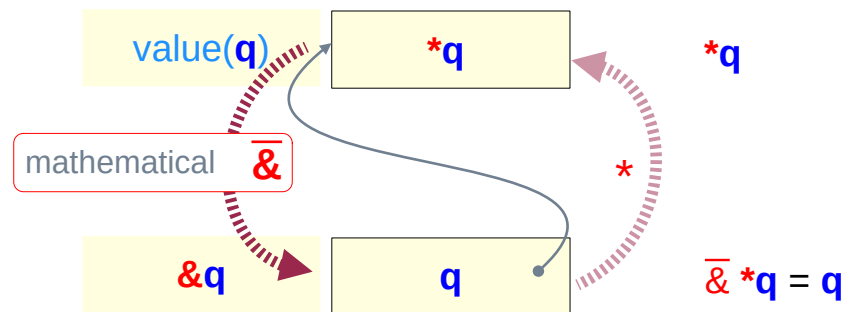


$$\bar{*} \&p = p$$

$\&p$

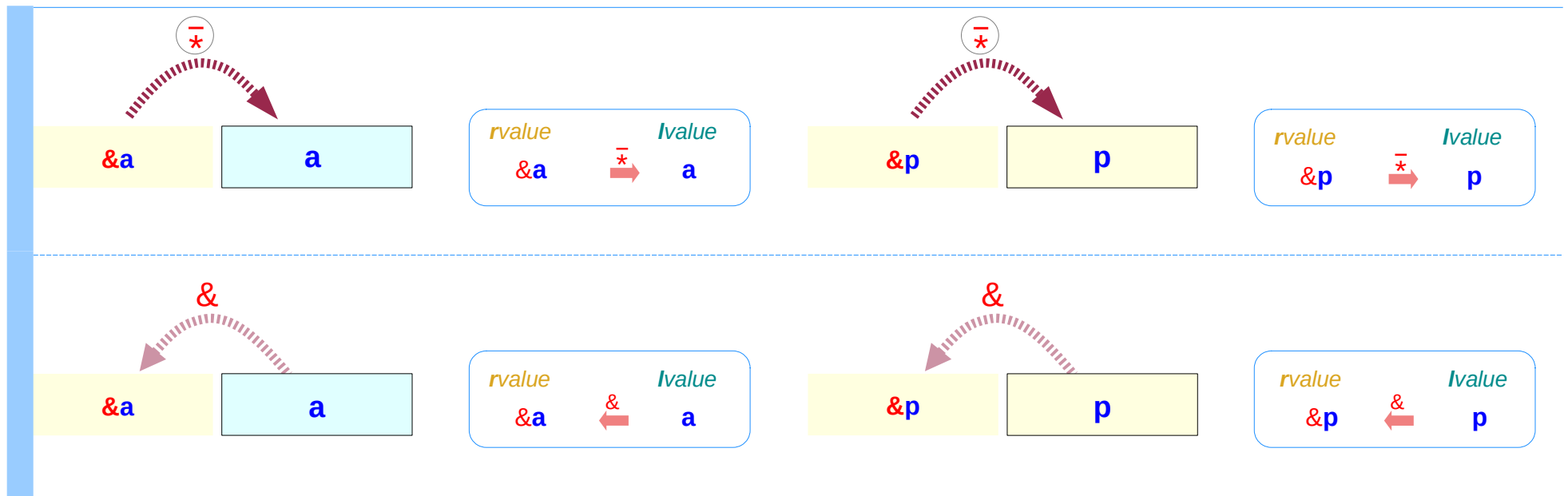


$$\bar{\&} *p = p$$



$$\bar{\&} *q = q$$

Inverse operators $\bar{*}$ and $\&$



$$\bar{*}\&a = a$$

$$\&\bar{*}\&a = \&a$$

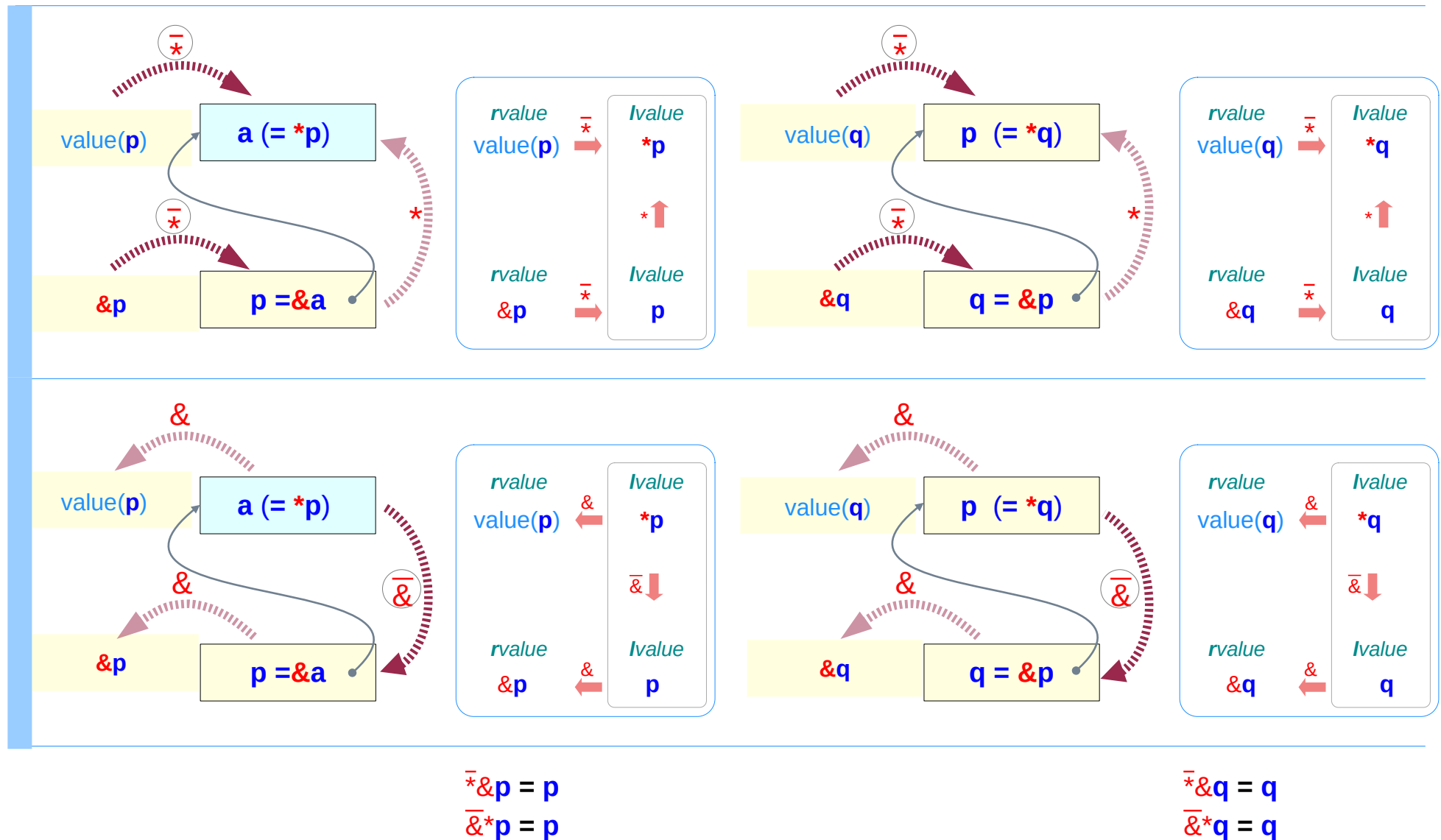
$\bar{*}$ and $\&$ are
inverse operators
to each other

$$\bar{*}\&p = p$$

$$\&\bar{*}\&p = \&p$$

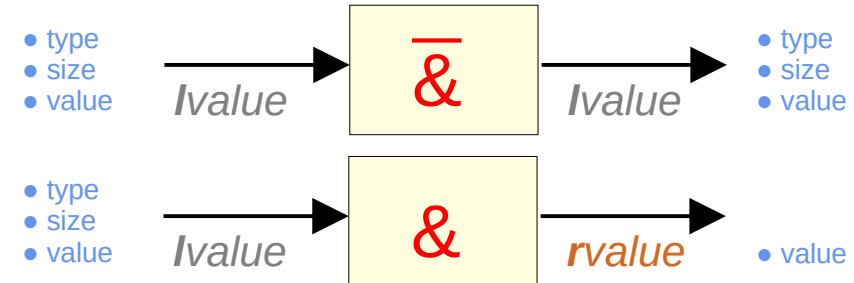
$\bar{*}$ and $\&$ are
inverse operators
to each other

Inverse operators $\bar{*}$ and $*$

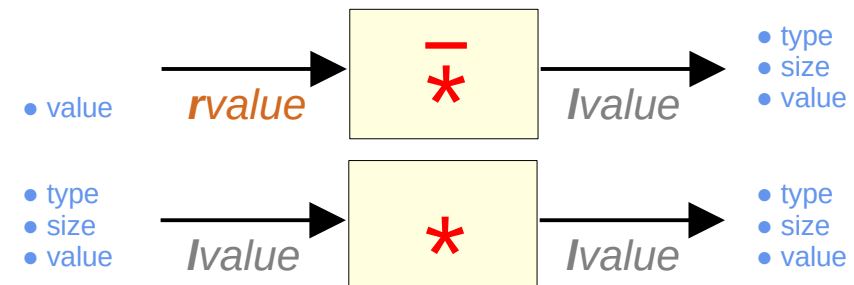
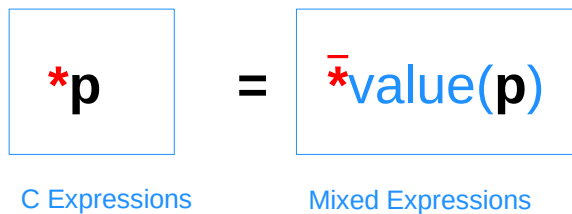


C operators and mathematical operators

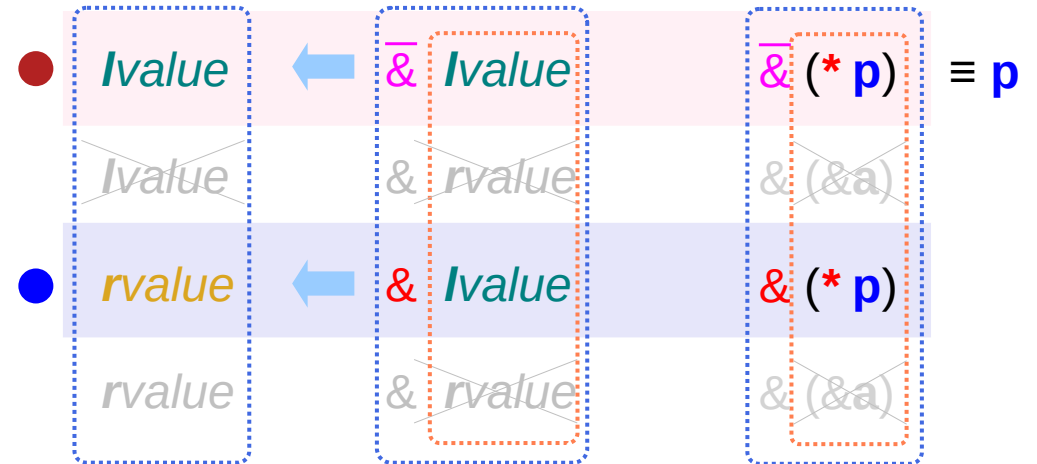
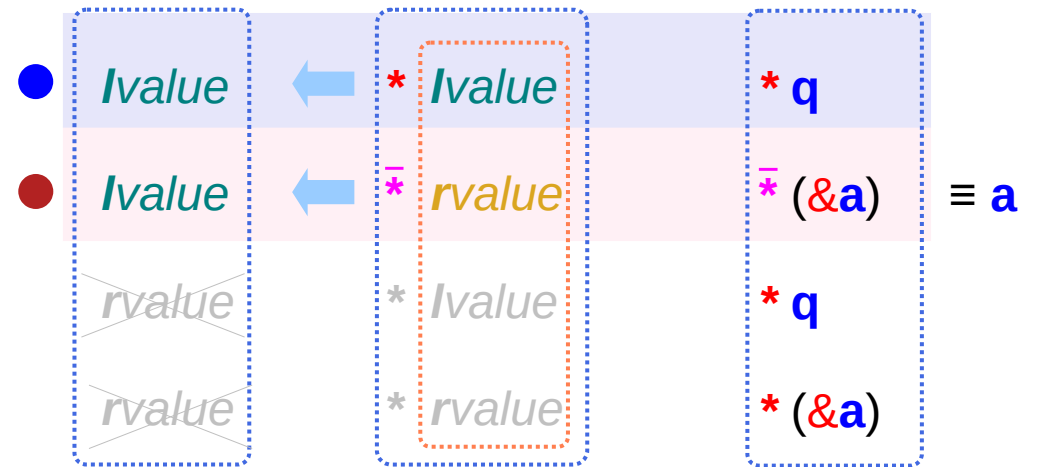
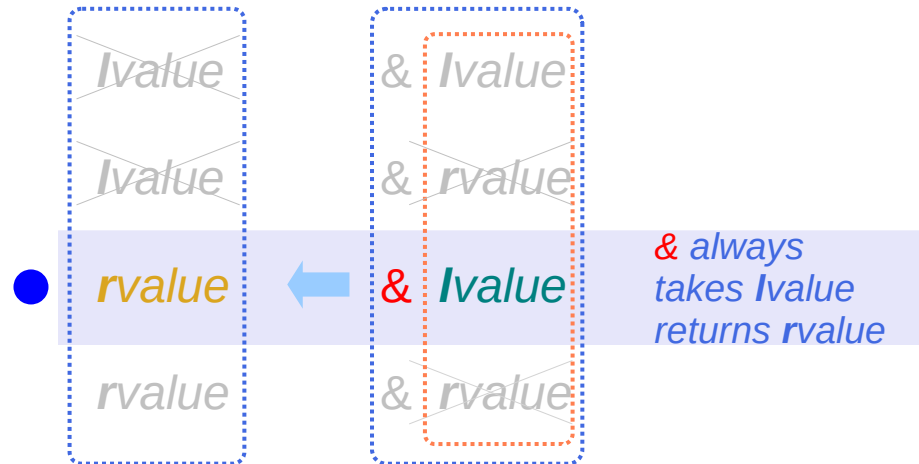
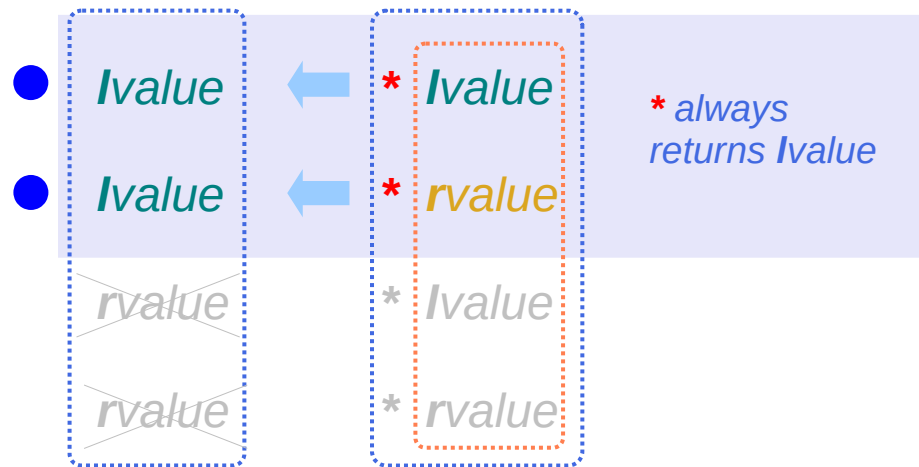
Address-of operation



De-reference operation



Ivalue and rvalue with * and & operators



a, p, q : Ivalues ... variables ... RW
 *p, *q, **q : Ivalues ... variables ... RW
 &a, &p, &q : rvalues ... constants ... RO

Examples of inverse operators

$\overline{*}\&a \rightarrow a$

$\overline{*}\&p \rightarrow p$

$\overline{*}\&q \rightarrow q$

$\overline{*}\text{value}(p) \rightarrow *p \rightarrow a$

$\overline{*}\text{value}(q) \rightarrow *q \rightarrow p$

$\overline{*}\text{value}(*q) \rightarrow **q \rightarrow *p \rightarrow a$

$\overline{\&}*p \rightarrow p$

$\overline{\&}*q \rightarrow q$

$\overline{\&}**q \rightarrow *q$

Extended Operators

$*\&a \rightarrow a$

$*\&p \rightarrow p$

$*\&q \rightarrow q$

$*p \rightarrow a$

$*q \rightarrow p$

$**q \rightarrow *p \rightarrow a$

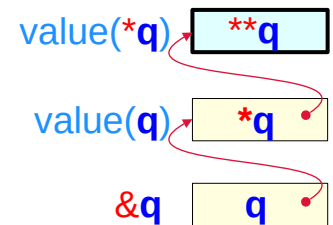
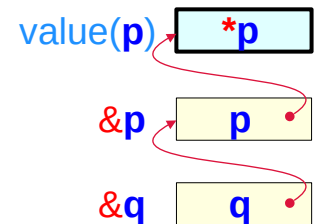
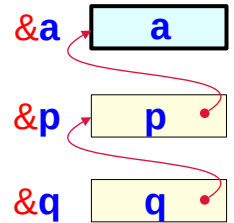
$\&*p \rightarrow \&a$

$\&*q \rightarrow \&p$

$\&**q \rightarrow \&a$

C Operators

```
int a;  
int * p = &a;  
int ** q = &p;
```



Operands of mathematical operators : $\bar{\&}$ and $\bar{*}$

$\bar{*}$ address value

$\bar{\&} \bar{*} \text{value}(\mathbf{P}) \longrightarrow \text{value}(\mathbf{P})$

$\bar{*} \bar{\&} \text{variable}$

$\bar{*} \bar{\&} \mathbf{X} \longrightarrow \mathbf{X}$

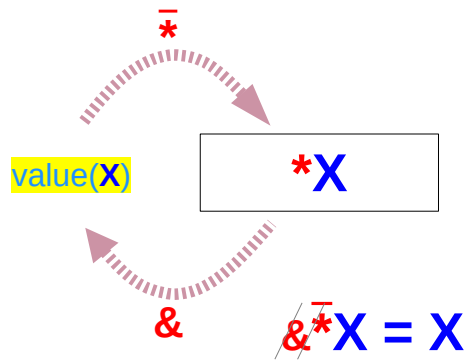
$\bar{\&} * \text{pointer}$
dereferenced pointer

$\bar{\&} * \mathbf{P} \longrightarrow \mathbf{P}$

$\bar{\&} \text{pointer}$
must be a dereferenced pointer
not considering
simultaneous pointing

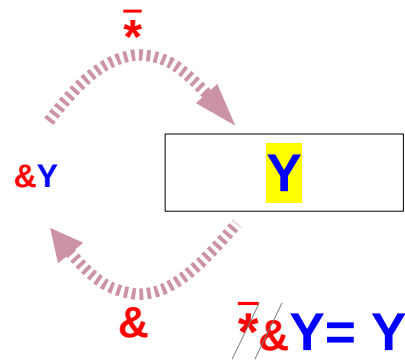
$* \bar{\&} \mathbf{Q} \longrightarrow \mathbf{Q}$
Q must be a dereferenced pointer
i.e, $\mathbf{Q} = *p \longrightarrow \bar{\&} \mathbf{Q} = p$

& and mathematical $\bar{*}$



$\&\bar{*}\text{value}(a) = \text{value}(a)$
 $\&\bar{*}\text{value}(p) = \text{value}(p)$
 $\&\bar{*}\text{value}(q) = \text{value}(q)$

$\&*a = \text{value}(a)$
 $\&*p = \text{value}(p)$
 $\&*q = \text{value}(q)$



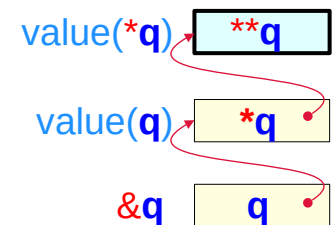
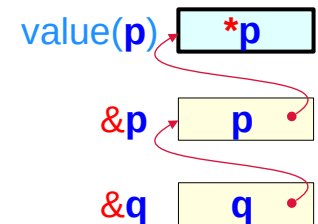
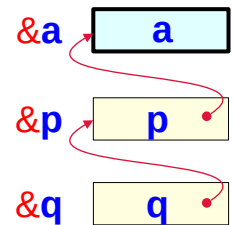
$\bar{*}\&a = a$
 $\bar{*}\&p = p$
 $\bar{*}\&q = q$

$\bar{*}\&a = a$
 $\bar{*}\&p = p$
 $\bar{*}\&q = q$

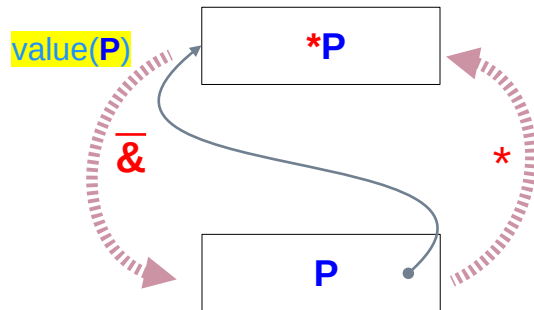
$*\&a = a$
 $*\&p = p$
 $*\&q = q$

C expressions

```
int a;  
int * p = &a;  
int ** q = &p;
```



* and mathematical $\bar{\&}$



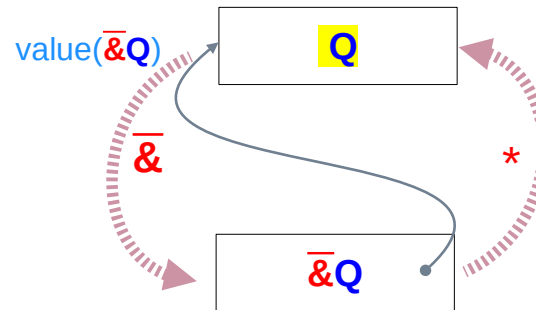
$$\bar{\&*P} = P$$

$$\bar{\&}*p = p$$

$$\bar{\&}*q = q$$

$$\&*p = \text{value}(p)$$

$$\&*q = \text{value}(q)$$



$$*\bar{\&Q} = Q$$

$$*\bar{\&}*p = *p$$

$$*\bar{\&}*q = *q$$

$$\bar{*}\bar{\&}*p = *p$$

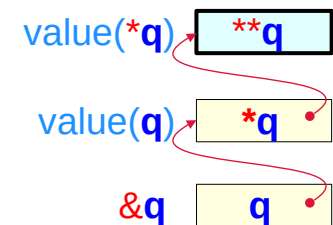
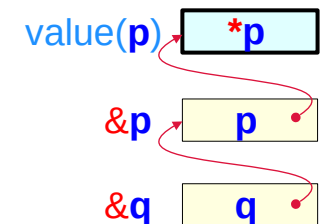
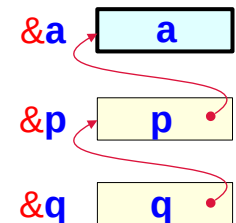
$$\bar{*}\bar{\&}*q = *q$$

$$\begin{aligned} * \&*p &= *p \\ * \&*q &= *q \end{aligned}$$

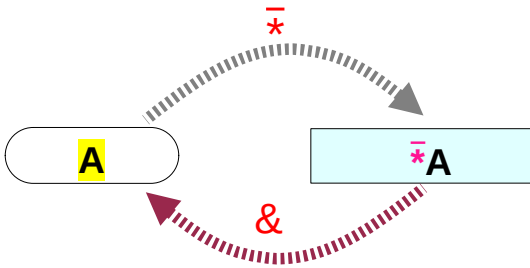
C operators

if not considering
simultaneous pointing

```
int a;
int * p = &a;
int ** q = &p;
```

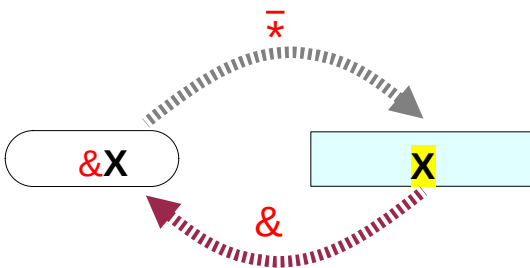


Requirements of address and data

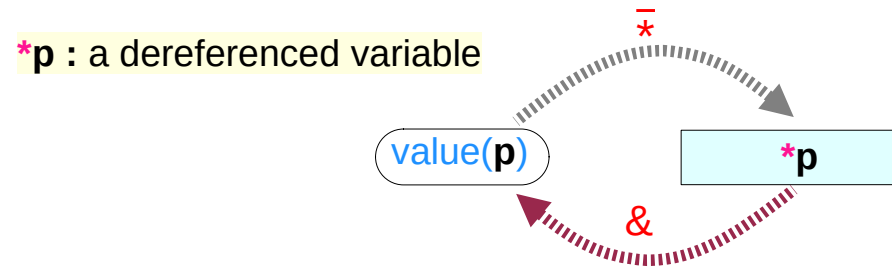


A : an address value - a number

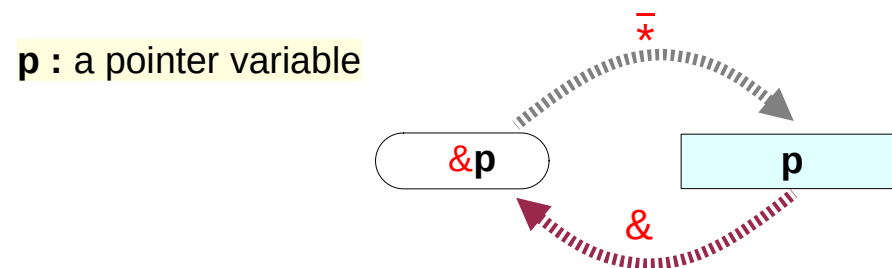
`value(p)`, `&p`, `&x`



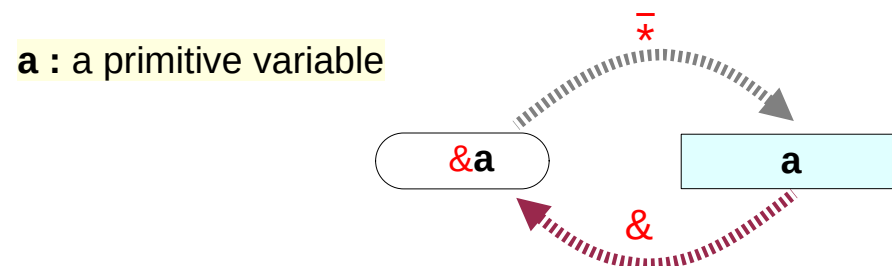
X : a variable `*p`, `p`, `a`



Pointer Dereferencing



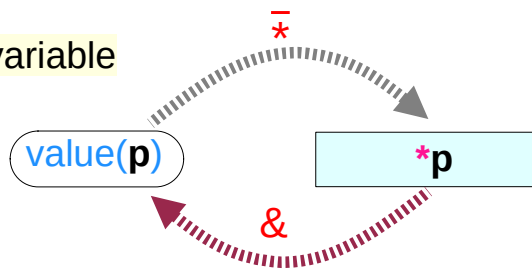
Pointer Data



Non-pointer Data

Inverse relations of $\&$ and $\bar{*}$ operators

$*p$: a dereferenced variable

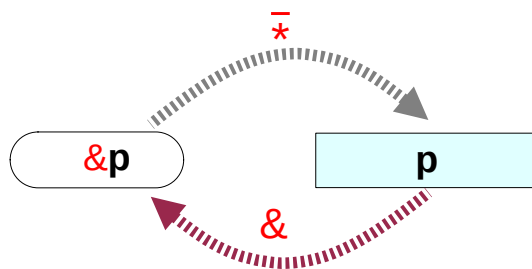


$$\bar{*} \text{value}(p) = *p$$

$$\&\bar{*} \text{value}(p) = \&*p = \text{value}(p)$$

$$\bar{*}\&\bar{*} \text{value}(p) = \bar{*}\&*p = \bar{*} \text{value}(p) = *p$$

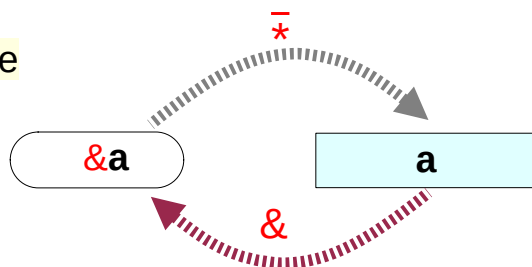
p : a pointer variable



$$\bar{*}\&p = p$$

$$\&\bar{*}\&p = \&p$$

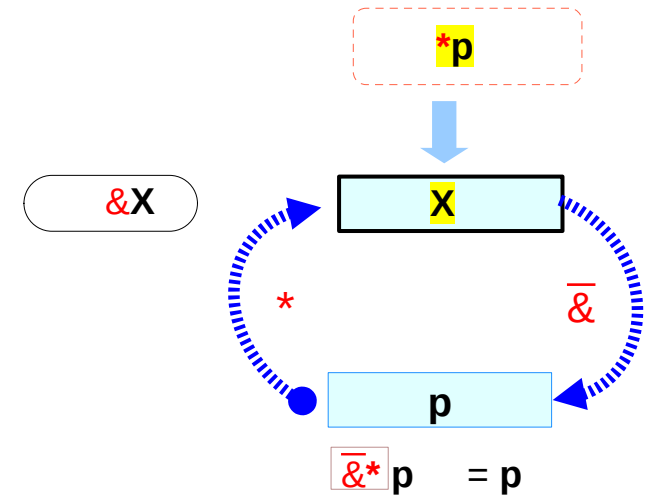
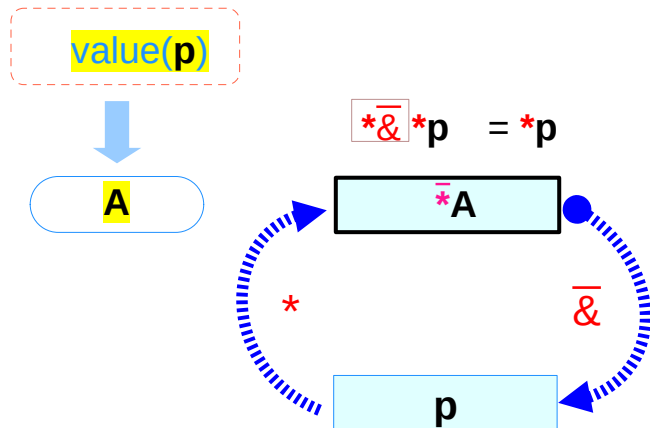
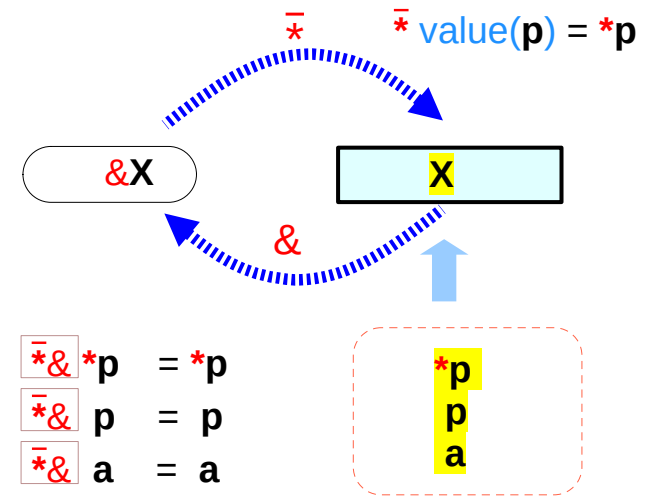
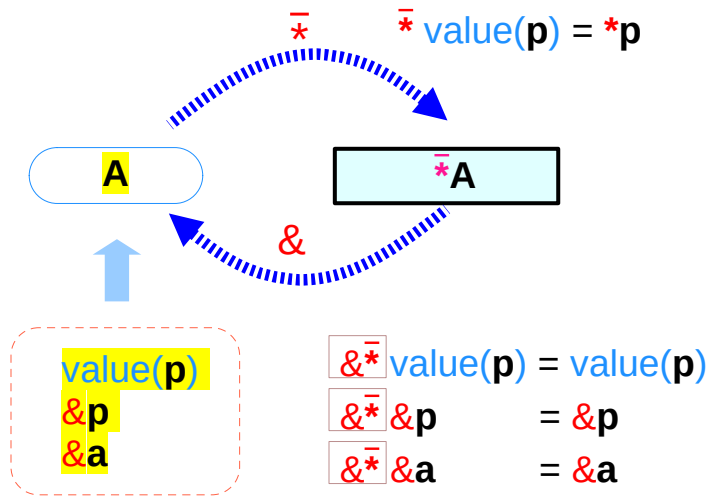
a : a primitive variable



$$\bar{*}\&a = a$$

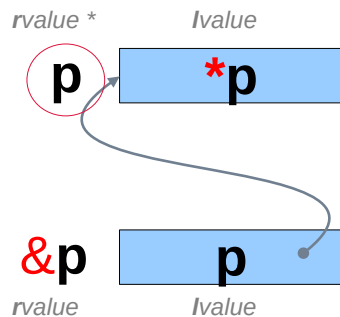
$$\&\bar{*}\&a = \&a$$

Requirements of pointed address and data

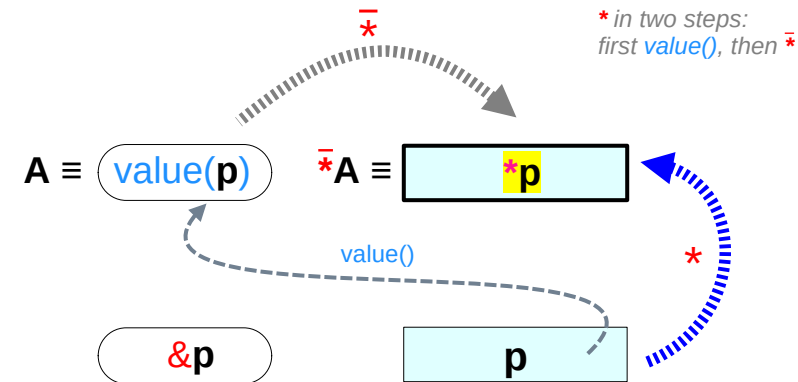


& and * operators in pointer de-referencing (1)

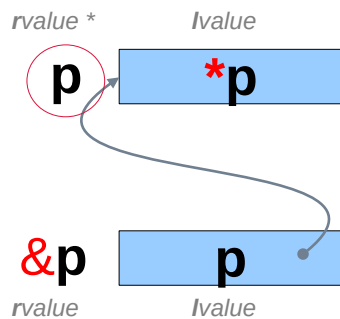
Two step De-reference * operation



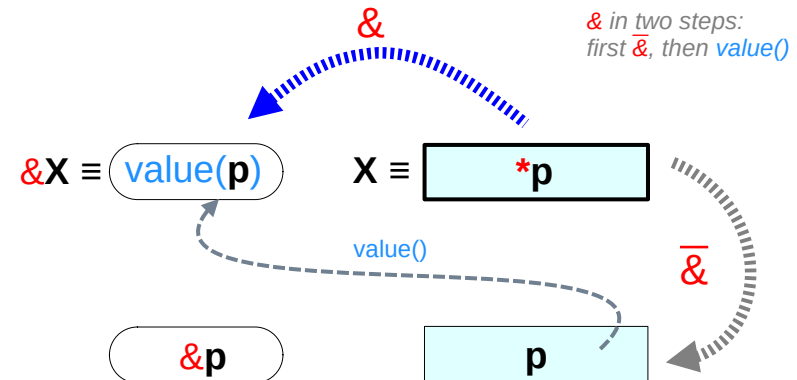
$$\begin{array}{lcl}
 \boxed{*p} & = & \boxed{\bar{*}\text{value}(p)} \\
 \text{C Expressions} & & \text{Mixed Expressions} \\
 \boxed{\bar{*}A} & = & \boxed{* \text{value}^{-1}(A)}
 \end{array}$$



Two step Address-of & operation

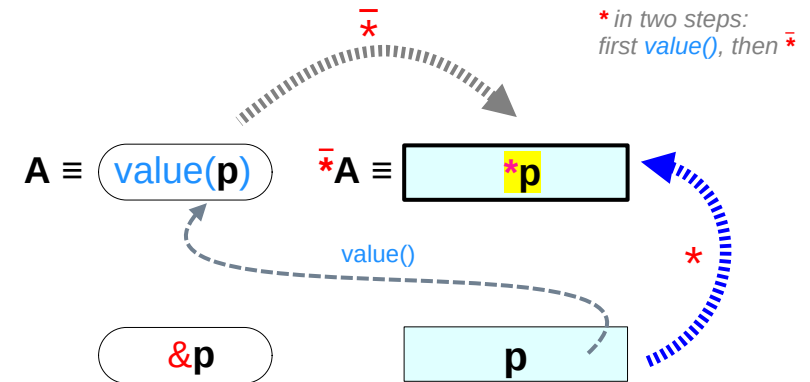
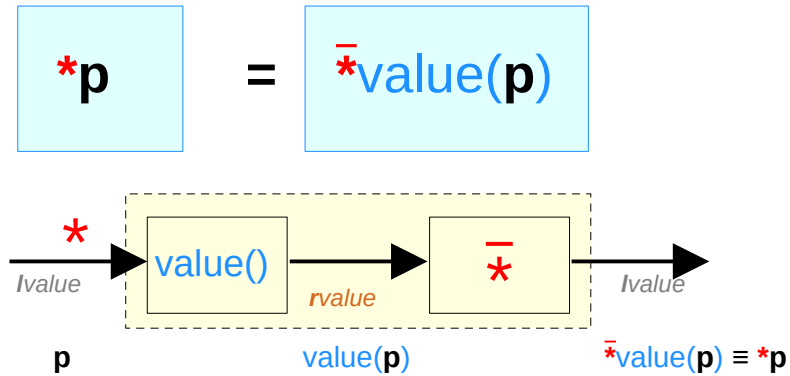


$$\begin{array}{lcl}
 \boxed{\&X} & = & \boxed{\text{value}(\bar{\&}X)} \\
 \text{C Expressions} & & \text{Mixed Expressions} \\
 \boxed{\bar{\&}X} & = & \boxed{\text{value}^{-1}(\&X)}
 \end{array}$$

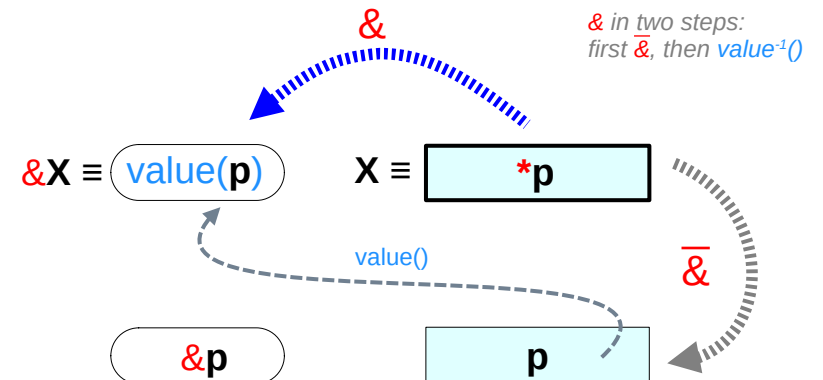
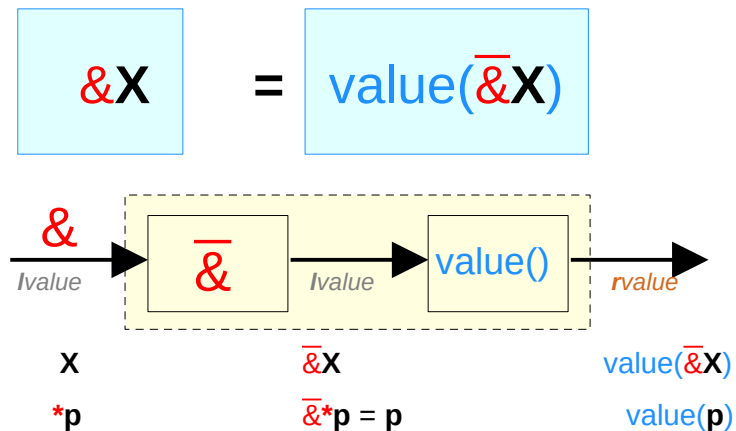


& and * operators in pointer de-referencing (2)

Two step De-reference * operation

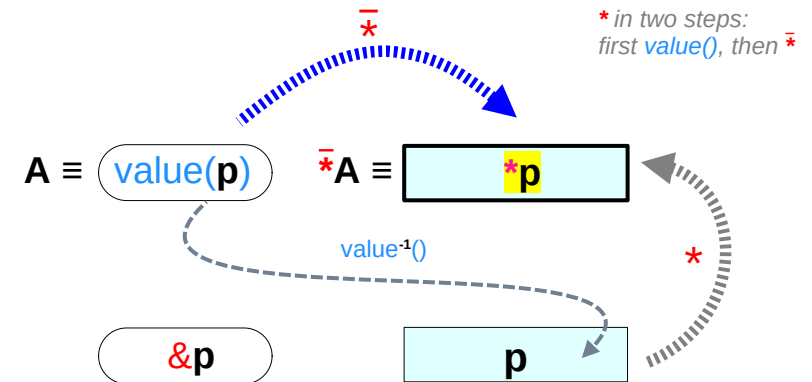
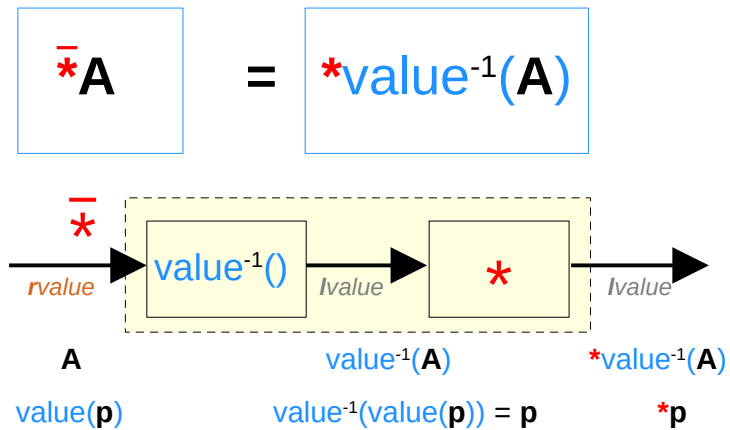


Two step Address-of & operation

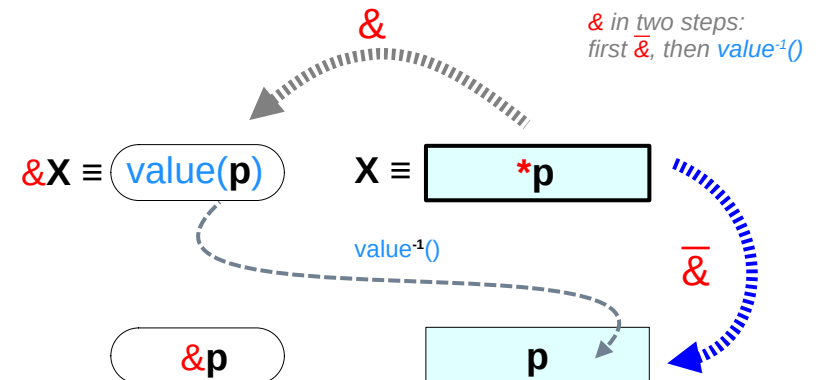
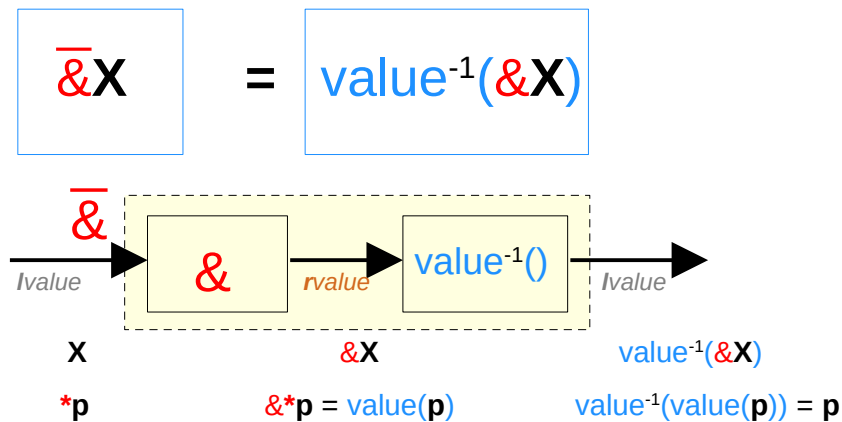


$\bar{*}$ and $\bar{*}$ operators in pointer de-referencing (3)

Two step De-reference $\bar{*}$ operation



Two step Address-of $\bar{\&}$ operation



Two step address-of & and $\bar{\&}$ operators

for $\&X$, X can be $*p$, p , a

$$\&X = \text{value}(\bar{\&X})$$

C Expressions

for $\bar{\&X} = p$, X must be $*p$

$$\bar{\&X} = \text{value}^{-1}(\&X)$$

$*p$ a dereferenced variable
 p a pointer variable
 a a primitive variable

$\&X$

$$\&*p = \&* \text{value}(p) = \text{value}(p)$$

$\&p$

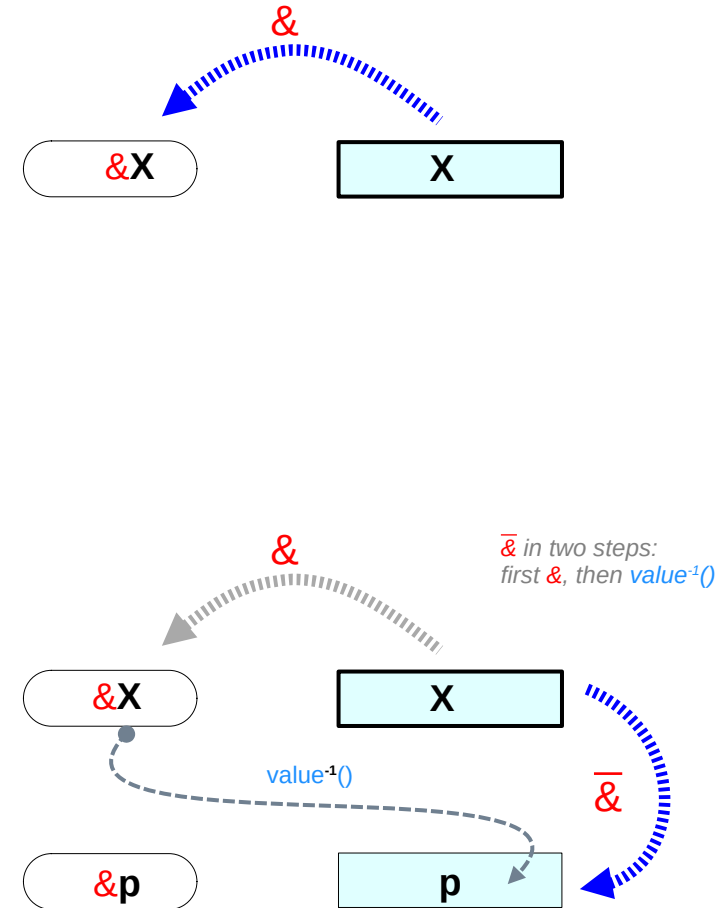
$\&a$

$*p$ a dereferenced variable
 ~~p a pointer variable~~
 ~~a a primitive variable~~

$$\bar{\&X} = p$$

$$\bar{\&*p} \equiv p$$

~~$\bar{\&p}$~~
 ~~$\bar{\&a}$~~



Two step de-reference $\bar{*}$ and $*$ operators

for $\bar{*}A$, A can be $\text{value}(p)$, $\&p$, $\&a$

$$\bar{*}A = * \text{value}^{-1}(A)$$

$\text{value}(p)$ value of a pointer variable
 $\&p$ address of a pointer variable
 $\&a$ address of a primitive variable

$$\begin{aligned} \bar{*}A \\ \bar{*} \text{value}(p) &= *p \\ \bar{*} \&p &= p \\ \bar{*} \&a &= a \end{aligned}$$

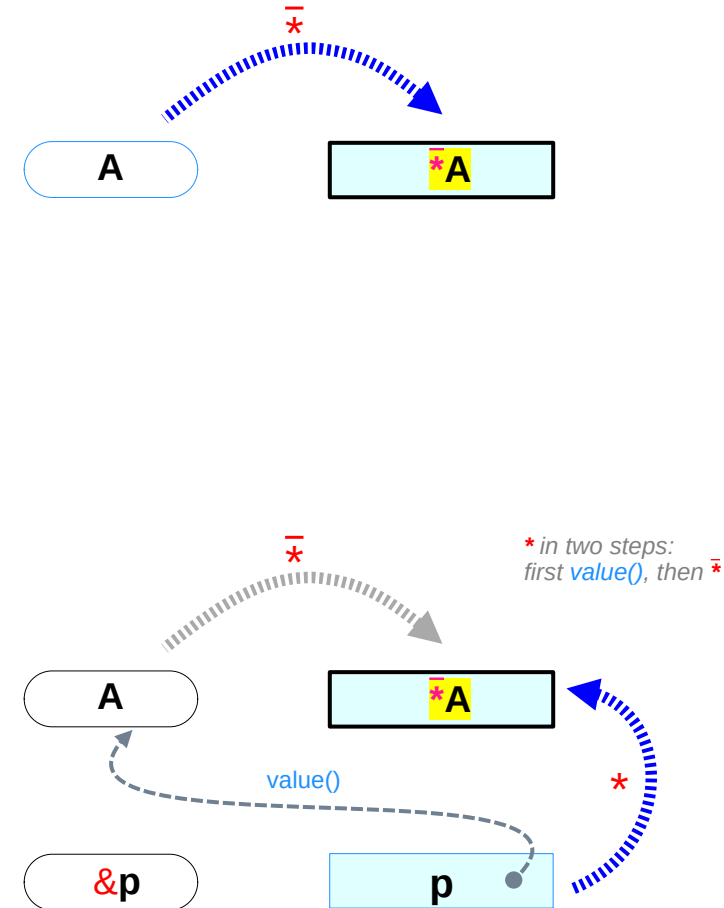
for $\bar{*}A = *p$, A must be $\text{value}(p)$

$$*p = \bar{*} \text{value}(p)$$

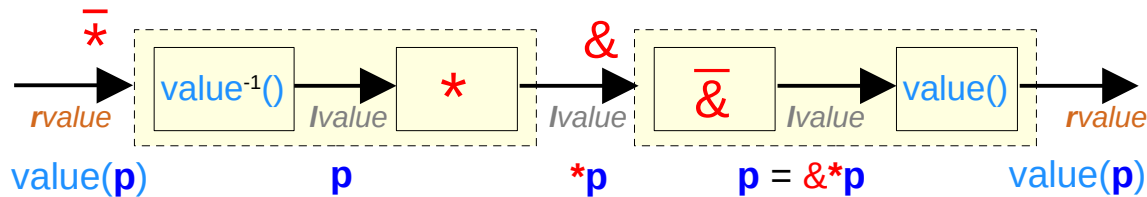
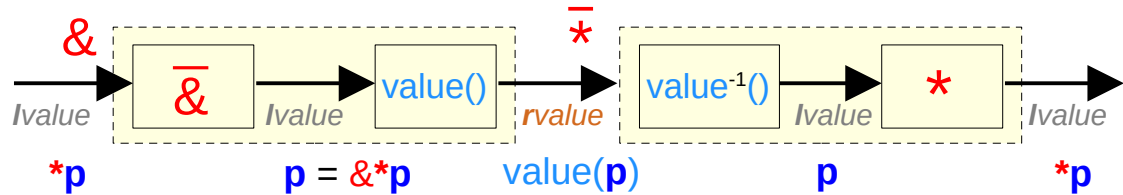
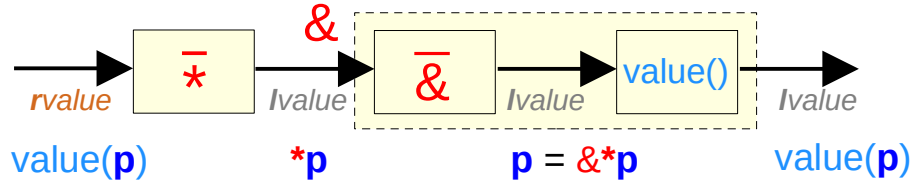
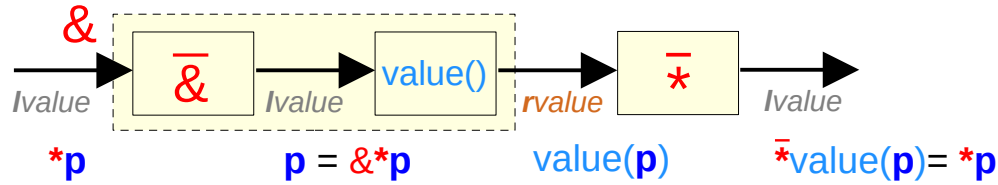
C Expressions

$\text{value}(p)$ value of a pointer variable
 ~~$\&p$ address of a pointer variable~~
 ~~$\&a$ address of a primitive variable~~

$$\begin{aligned} *p \\ *p &\equiv \bar{*} \text{value}(p) \\ \cancel{* \&p} &= p \\ \cancel{* \&a} &= a \end{aligned}$$



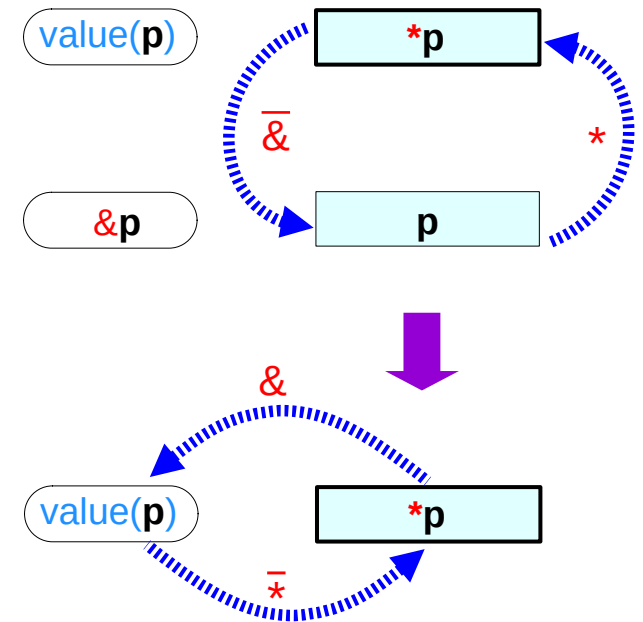
Inverse operators : $\&$ and $\bar{*}$ in pointer de-referencing



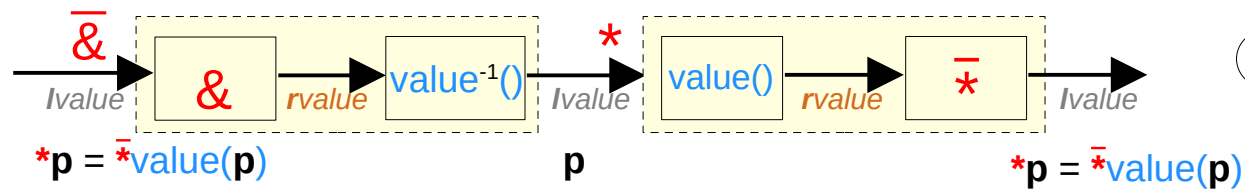
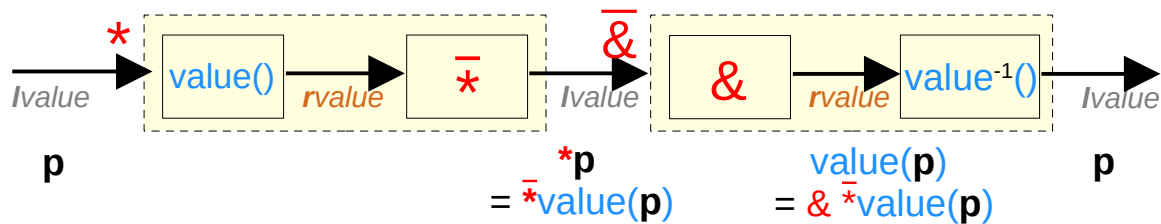
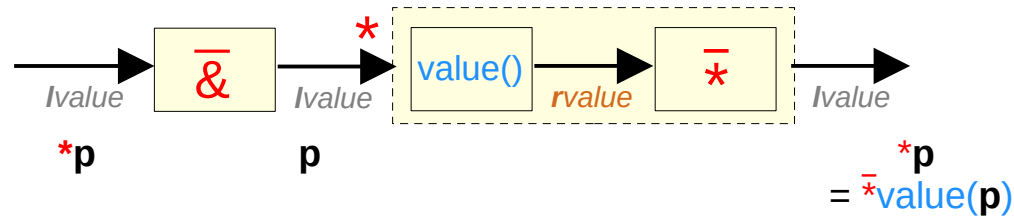
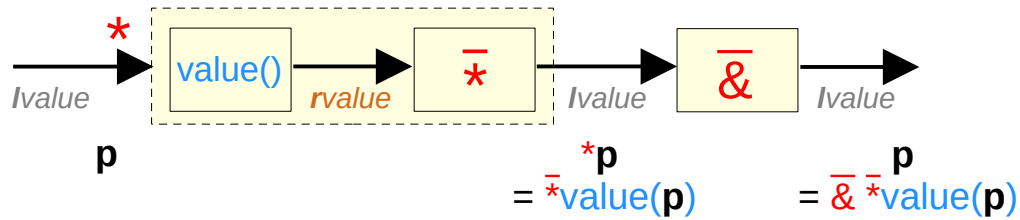
$$\bar{*} value(p) = *p$$

$$\&*value(p) = \&*p = value(p)$$

$$\bar{*}\&*value(p) = \bar{*}\&*p = \bar{*}value(p) = *p$$



Inverse operators of & and *



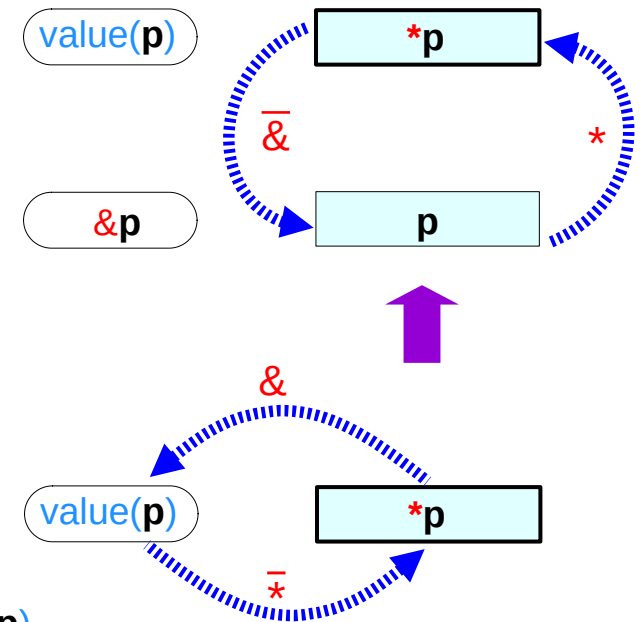
$$\bar{*} value(p) = *p$$

$$\bar{\&} * p = p$$

$$* \bar{\&} * p = *p$$

$$p = \bar{\&} \bar{*} value(p)$$

$$value(p) = \bar{\&} \bar{*} value(p)$$



Inverse operators in pointer referencing (1)

$$\begin{aligned}\overline{*}&\&*p &= *p \\ \overline{*}&\&p &= p \\ \overline{*}&\&a &= a\end{aligned}$$

$$\begin{aligned}\&\overline{*} \text{value}(p) &= \text{value}(p) \\ \&\overline{*} \&p &= \&p \\ \&\overline{*} \&a &= \&a\end{aligned}$$

$\overline{*}&$

$*p$
 p
 a

$\overline{\&*}$

p

$\&\overline{*}$

$\text{value}(p)$
 $\&p$
 $\&x$

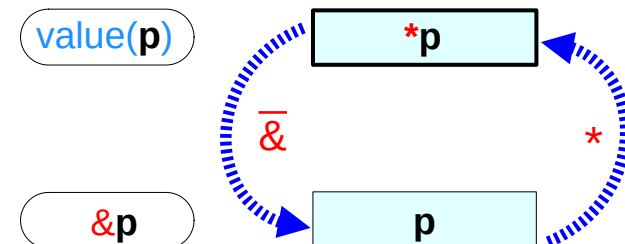
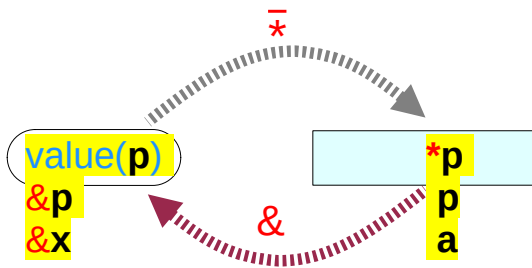
$*\overline{\&}$

$*p$

$$\overline{\&*} p = p$$

$$\&\overline{*} *p = *p$$

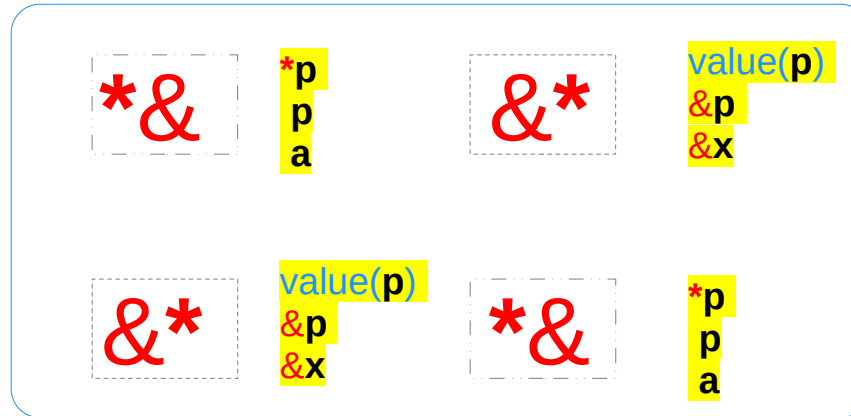
Math expressions



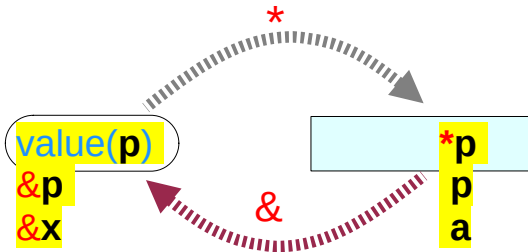
Inverse operators in pointer referencing (2)

`*&*p = *p`
`*&p = p`
`*&a = a`

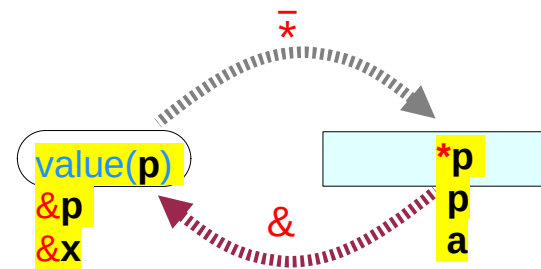
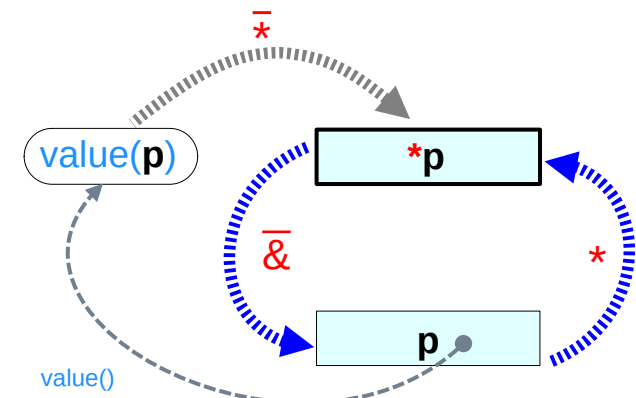
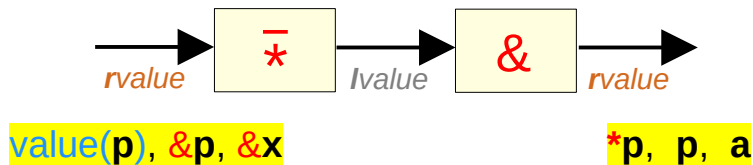
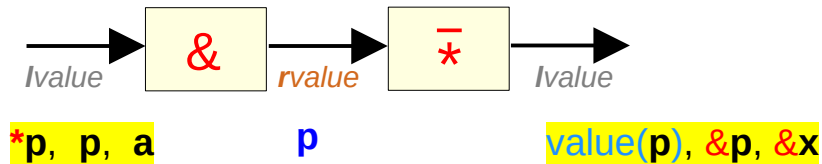
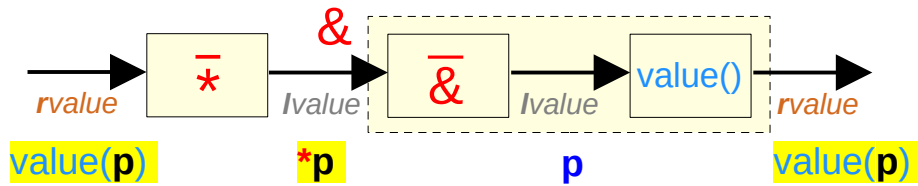
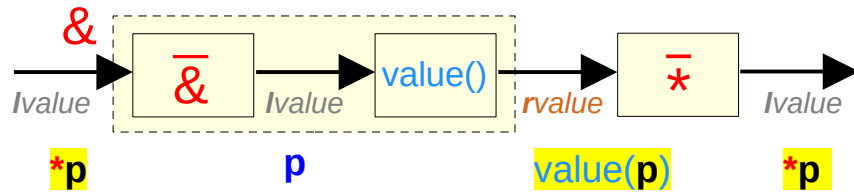
`&*value(p) = value(p)`
`&*&p = &p`
`&*&a = &a`



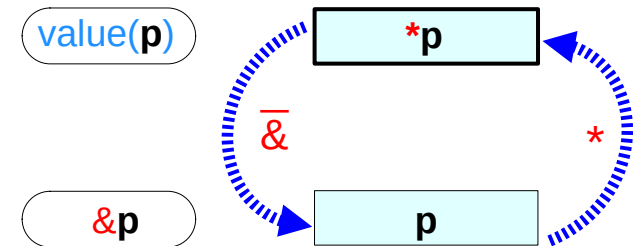
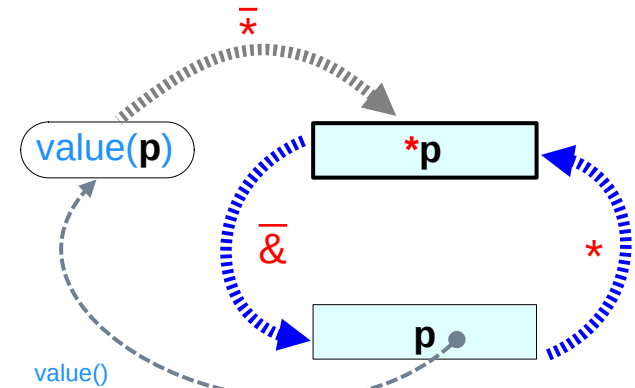
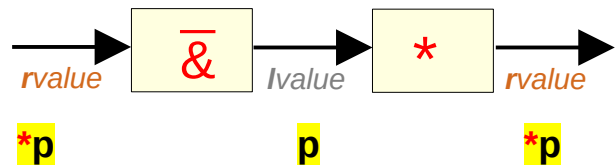
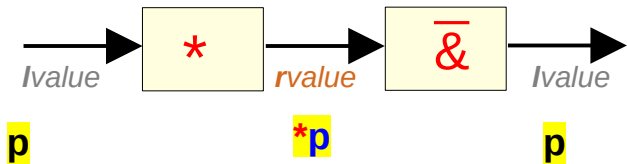
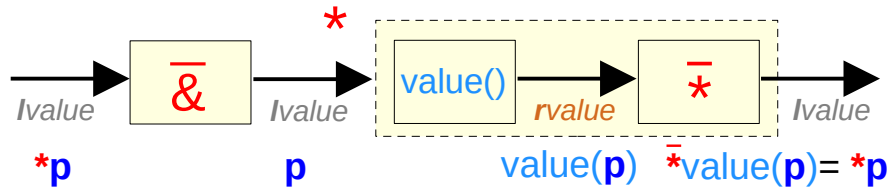
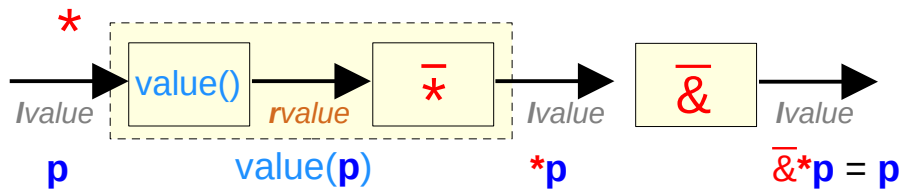
C expressions



Inverse operators in pointer referencing (3)



Inverse operators in pointer referencing (3)



References

- [1] Essential C, Nick Parlante
- [2] Efficient C Programming, Mark A. Weiss
- [3] C A Reference Manual, Samuel P. Harbison & Guy L. Steele Jr.
- [4] C Language Express, I. K. Chun